FORM-PTO-1390 (Rev. 12-29-99) U.S. DEPARTMENT OF COMMERCE PATENT AND TRADEMARK OFFICE

TRANSMITTAL LETTER TO THE UNITED STATES DESIGNATED/ELECTED OFFICE (DO/EO/US) CONCERNING A FILING UNDER 35 U.S.C. 371

ATTORNEY'S DOCKET NUMBER: 032326-133

U.S. APPLICATION NO. (If known, see 37 C.F.R. 1.5): Unassigned 09/807615

INTERNATIONAL APPLICATION NO.	INTERNATIONAL FILING DATE	PRIORITY DATE CLAIMED
PCT/FR99/02199	15 September 1999	16 October 1998

TITLE OF INVENTION: COUNTERMEASURE METHOD IN AN ELECTRONIC COMPONENT USING A SECRET KEY CRYPTOGRAPHIC ALGORITHM

APPLICANT(S) FOR DO/EO/US: Christophe CLAVIER and Olivier BENOIT

Applicant herewith submits to the United States Designated/Elected Office (DO/EO/US) the following items and other information:

1. ☒ This is a **FIRST** submission of items concerning a filing under 35 U.S.C. 371.
2. ☐ This is a **SECOND** or **SUBSEQUENT** submission of items concerning a filing under 35 U.S.C. 371.
3. ☒ This is an express request to begin national examination procedures (35 U.S.C. 371(f)) at any time rather than delay examination until the expiration of the applicable time limit set in 35 U.S.C. 371(b) and the PCT Articles 22 and 39(1).
4. ☒ A proper Demand for International Preliminary Examination was made by the 19th month from the earliest claimed priority date.
5. ☒ A copy of the International Application as filed (35 U.S.C. 371(c)(2))
   a. ☒ is transmitted herewith (required only if not transmitted by the International Bureau).
   b. ☒ has been transmitted by the International Bureau.
   c. ☐ is not required, as the application was filed in the United States Receiving Office (RO/US)
6. ☒ A translation of the International Application into English (35 U.S.C. 371(c)(2)).
7. ☒ Amendments to the claims of the International Application under PCT Article 19 (35 U.S.C. 371(c)(3))
   a. ☐ are transmitted herewith (required only if not transmitted by the International Bureau).
   b. ☐ have been transmitted by the International Bureau.
   c. ☐ have not been made; however, the time limit for making such amendments has NOT expired.
   d. ☒ have not been made and will not be made.
8. ☐ A translation of the amendments to the claims under PCT Article 19 (35 U.S.C. 371(c)(3)).
9. ☐ An oath or declaration of the inventor(s) (35 U.S.C. 371(c)(4)).
10. ☐ A translation of the annexes to the International Preliminary Examination Report under PCT Article 36 (35 U.S.C. 371(c)(5)).

Items 11. to 16. below concern other document(s) or information included:

11. ☒ An Information Disclosure Statement under 37 CFR 1.97 and 1.98.
12. ☐ An assignment document for recording. A separate cover sheet in compliance with 37 CFR 3.28 and 3.31 is included.
13. ☒ A FIRST preliminary amendment.
    ☐ A SECOND or SUBSEQUENT preliminary amendment.
14. ☐ A substitute specification.
15. ☐ A change of power of attorney and/or address letter.
16. ☐ Other items or information:

U.S. APPLICATION NO. (If known, see 37 C.F.R. 1.5)	INTERNATIONAL APPLICATION NO.	ATTORNEY'S DOCKET NUMBER
Unassigned 09/807615	PCT/FR99/02199	032326-133

17. ☒ The following fees are submitted:	CALCULATIONS	PTO USE ONLY
Basic National Fee (37 CFR 1.492(a)(1)-(5)):		
Neither international preliminary examination fee (37 CFR 1.482) nor international search fee (37 CFR 1.445(a)(2)) paid to USPTO and International Search Report not prepared by the EPO or JPO $1,000.00 (960)		
International preliminary examination fee (37 CFR 1.482) not paid to USPTO but International Search Report prepared by the EPO or JPO $860.00 (970)		
International preliminary examination fee (37 CFR 1.482) not paid to USPTO but international search fee (37 CFR 1.445(a)(2)) paid to USPTO $710.00 (958)		
International preliminary examination fee paid to USPTO (37 CFR 1.482) but all claims did not satisfy provisions of PCT Article 33(1)-(4) $690.00 (956)		
International preliminary examination fee paid to USPTO (37 CFR 1.482) and all claims satisfied provisions of PCT Article 33(1)-(4) $100.00 (962)		
ENTER APPROPRIATE BASIC FEE AMOUNT =	$ 860.00	
Surcharge of **$130.00 (154)** for furnishing the oath or declaration later than 20 ☐ 30 ☐ months from the earliest claimed priority date (37 CFR 1.492(e)).	$ -0-	

Claims	Number Filed	Number Extra	Rate		
Total Claims	15 -20 =	-0-	X$18.00 (966)	$ -0-	
Independent Claims	2 -3 =	-0-	X$80.00 (964)	$ -0-	
Multiple dependent claim(s) (if applicable)			+$270.00 (968)	$ -0-	
TOTAL OF ABOVE CALCULATIONS =				$ 860.00	
Reduction for 1/2 for filing by small entity, if applicable (see below).				$ -0-	-
SUBTOTAL =				$ 860.00	
Processing fee of **$130.00 (156)** for furnishing the English translation later than 20 ☐ 30 ☐ months from the earliest claimed priority date (37 CFR 1.492(f)). +				$ -0-	
TOTAL NATIONAL FEE =				$ -0-	
Fee for recording the enclosed assignment (37 CFR 1.21(h)). The assignment must be accompanied by an appropriate cover sheet (37 CFR 3.28, 3.31). **$40.00 (581)** per property +				$ -0-	
TOTAL FEES ENCLOSED =				$ 860.00	
				Amount to be: refunded	$
				charged	$

a. ☐ Small entity status is hereby claimed.

b. ☒ A check in the amount of $ 860.00 to cover the above fees is enclosed.

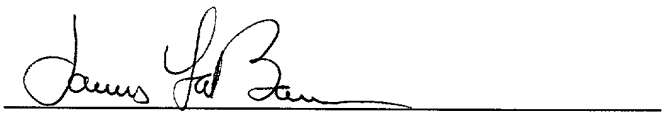
c. ☐ Please charge my Deposit Account No. 02-4800 in the amount of $_____ to cover the above fees. A duplicate copy of this sheet is enclosed.

d. ☐ The Commissioner is hereby authorized to charge any additional fees which may be required, or credit any overpayment to Deposit Account No. 02-4800. A duplicate copy of this sheet is enclosed.

NOTE: Where an appropriate time limit under 37 CFR 1.494 or 1.495 has not been met, a petition to revive (37 CFR 1.137(a) or (b)) must be filed and granted to restore the application to pending status.

SEND ALL CORRESPONDENCE TO:

James A. LaBarre
BURNS, DOANE, SWECKER & MATHIS, L.L.P.
P.O. Box 1404
Alexandria, Virginia 22313-1404
(703) 836-6620



SIGNATURE

James A. LaBarre
NAME

28,632
REGISTRATION NUMBER

Patent
Attorney's Docket No. 032326-133

IN THE UNITED STATES PATENT AND TRADEMARK OFFICE

In re Patent Application of	)	
	)	
Christophe CLAVIER et al	)	Group Art Unit: Unassigned
	)	
Application No.: Unassigned	)	Examiner: Unassigned
	)	
Filed: April 16, 2001	)	
	)	
For: COUNTERMEASURE METHOD IN AN ELECTRONIC COMPONENT USING A SECRET KEY CRYPTOGRAPHIC ALGORITHM	)	

PRELIMINARY AMENDMENT

Assistant Commissioner for Patents
Washington, D.C. 20231

Sir:

Prior to examination and the calculation of filing fees, kindly amend the above-identified application as follows:

IN THE SPECIFICATION:

Page 1, immediately following the title appearing on lines 1 and 2, insert the following:

--This disclosure is based upon, and claims priority from French Application No. 98/12990, filed on October 16, 1998 and International Application No. PCT/FR99/02199, filed September 15, 1999, which was published on April 27, 2000 in a language other than English, the contents of which are incorporated herein by reference.

Background of the Invention--

Page 8, between lines 26 and 27, insert the following heading:

--**Summary of the Invention** --.

Page 11, between lines 16 and 17, insert the following heading:

--**Brief Description of the Drawings**--.

Page 13, between lines 6 and 7, insert the following heading:

--**Detailed Description**--.

Kindly replace the paragraph beginning at page 17, line 7, with the following:

The data item 1 of a given round is a data item derived from the data item g of the previous round, since it corresponds to a permutation of the bits of the word g, certain bits of the word g also being duplicated.

Kindly replace the paragraph beginning at page 28, line 20, with the following:

It is clear from this flow diagram that complemented data are obtained (the complementation being denoted by a bar above the data item) for all the critical instructions of these rounds. And the data L3 and R3 at the output of the third round are not complemented. The execution of the algorithm can be continued by passing to the round T4, at which the first means TC_0 are applied according to the normal execution of the algorithm.

IN THE CLAIMS:

Cancel claims 9 and 10.

Kindly replace claims 1-8, as follows.

1. (Amended) A countermeasure method in an electronic component using a secret key cryptographic algorithm for calculating an encoded message from an input message, of the type in which sixteen calculation rounds are employed where each round supplies an output data item from an input data item, and the output data item is manipulated by critical instructions in at least the first three and last three rounds, said method including the following steps:

forming a group comprising at least the first three rounds and another group comprising at least the last three rounds,

in each of these groups selectively applying a first sequence that uses a first manipulating means for said critical instructions in each round or a second sequence that uses other manipulating means for said critical instructions at least in certain rounds, said first and second sequences being such that they supply the same result at the output from the last round in each group for the same given input message, and

selecting the sequence to be executed in the groups as a function of a statistical half probability distribution, in order to make the data manipulated by said critical instructions unpredictable.

2. (Amended) A countermeasure method according to Claim 1, wherein each of said manipulating means produces output data in accordance with input data, and

wherein said other manipulating means are such that they complement at least one or both of the input and/or output data of the first manipulating means.

3. (Amended) A countermeasure method according to Claim 2, wherein said second sequence comprises, for one or more rounds, an additional complementation operation at the input or output of the manipulating means used, and wherein said first sequence includes an additional operation of identical copying that corresponds to each additional complementation operation in said second sequence.

4. (Amended) A countermeasure method according to claim 1, wherein four groups of each of four successive rounds are formed, and wherein said first sequence is applied to each group and said second sequence is applied to at least the first group and the last group.

5. (Amended) A countermeasure method according to Claim 4, wherein second sequence is applied to each of the groups.

6. (Amended) A countermeasure method according to claim 1, wherein the first group is formed by the first three rounds and the last group is formed by the last three rounds.

7. (Amended) A countermeasure method according to claim 1, wherein the step of selecting the sequence to be executed is made at the start of execution of the algorithm by drawing a random value.

8. (Amended) A countermeasure method according to claim 1, wherein said manipulating means are tables of constants.

Add the following new claims:

--11. (New) An electronic security component have a countermeasure against attacks on a secret key cryptography technique in which data is manipulated by critical instructions, said component comprising:

a program memory having stored therein a plurality of manipulating means for use during said critical instructions, said manipulating means having complementary input and/or output data relative to one another; and

means for generating a random value that designates at least one of said manipulating means to be employed during a given execution of said cryptography technique.

12. (New) The electronic security component of claim 11, wherein said plurality of manipulating means each comprise a table of constants.

13. (New) The electronic security component of claim 11, wherein said cryptography technique comprises a DES algorithm that is executed in multiple rounds.

14. (New) The electronic security component of claim 13 wherein said random value has a first state which designates a manipulating means that is to be employed during all of the rounds of said algorithm, and a second state which designates at least two other manipulating means that are to be employed during different respective rounds of said algorithm.

15. (New) The electronic security component of claim 11, wherein said component is a chip card.--

REMARKS

Entry of the foregoing amendment is respectfully requested. This amendment is intended to place the claims in a more conventional format and eliminate the multiple dependency of the claims.

Respectfully submitted,

BURNS, DOANE, SWECKER & MATHIS, L.L.P.

By: _____

James A. LaBarre
Registration No. 28,632

P.O. Box 1404
Alexandria, Virginia 22313-1404
(703) 836-6620

Date: April 16, 2001

Attachment to Preliminary Amendment dated April 16, 2001

Marked-up Copy

Kindly replace the paragraph beginning at page 17, line 7, with the following:

The data item 1 of [the first] a given round is a data item derived from the data item g of the [first] previous round, since it corresponds to a permutation of the bits of the word g, certain bits of the word g also being duplicated.

Kindly replace the paragraph beginning at page 28, line 20, with the following:

It is clear from this flow diagram that complemented data are obtained (the complementation being denoted by a bar above the data item) for all the critical instructions of these rounds. And the data L3 and R3 at the output of the third round are not complemented. The execution of the algorithm can be continued by passing to the round T4, at which the first means [TC_3] $\underline{TC_0}$ are applied according to the normal execution of the algorithm.

Attachment to Preliminary Amendment dated April 16, 2001

Marked-up Claims 1-8

1. (Amended) A countermeasure method in an electronic component using a secret key cryptographic algorithm [(K)] for calculating an encoded message [(C)] from an input message [(M)], [the use of the algorithm comprising] of the type in which sixteen calculation rounds [(T1, ..., T16), each round using first means (TC_0) for supplying] are employed where each round supplies an output data item from an input data item, and the output data item [and/or the derived data being] is manipulated by critical instructions in at least the first three [(T1, T2, T3)] and last three [(T14, T15, T16)] rounds, [characterised in that] said method including the following steps:

forming a group [(G1) is formed] comprising at least the first three rounds and another group [(G4)] comprising at least the last three rounds, [and]

in [that with] each of these groups [(G1 and G4) there is associated] selectively applying a first sequence [(SEQA) using the] that uses a first manipulating means [(TC_0)] for said critical instructions in each round [and] or a second sequence [(SEQB) using] that uses other manipulating means [(TC_1, TC_2, TC_3) in] for said critical instructions at least in certain rounds [at least], [the] said first and second sequences being such that they supply the same result at the output from the last round in each group for the same given input message [(M)], and

[the choice of] selecting the sequence to be executed in the groups [concerned being] as a function of a statistical half probability distribution, in order to make [all] the data manipulated by [the] said critical instructions unpredictable.

Attachment to Preliminary Amendment dated April 16, 2001

Marked-up Claims 1-8

2. (Amended) A countermeasure method according to Claim 1, [characterised in that the other] wherein each of said manipulating means produces output data in accordance with input data, and wherein said other manipulating means are such that they complement at least one [or other] or both of the input [(E)] and/or output [(S)] data of the first manipulating means.

3. (Amended) A countermeasure method according to Claim 2, [characterised in that the] wherein said second sequence [(SEQB)] comprises, for one or more rounds, an additional complementation operation [(CP)] at the input or output of the manipulating means used, and [in that, to each additional complementation] wherein said first sequence includes an additional operation of identical copying [(ID) in the first sequence (SEQA)] that corresponds to each additional complementation operation in said second sequence.

4. (Amended) A countermeasure method according to [any one of the preceding claims, characterised in that] claim 1, wherein four groups [(G1, ... G4)] of each of four successive rounds [(T1, ... T4)] are formed, [in that the] and wherein said first sequence [(SEQA) is associated with] is applied to each group and [in that the] said second sequence [(SEQB) is associated] is applied to at least [with] the first group [(G1)] and the last group [(G4)].

Attachment to Preliminary Amendment dated April 16, 2001

Marked-up Claims 1-8

5. (Amended) A countermeasure method according to Claim 4, [characterised in that the] wherein second sequence [(SEQB) is associated with] is applied to each of the groups [(G1, ... G4)].

6. (Amended) A countermeasure method according to [any one of Claims 1 to 3, characterised in that] claim 1, wherein the first group [(G1)] is formed by the first three rounds [(T1, T2, T3)] and [in that] the last group is formed by the last three rounds [(T14, T15, T16)].

7. (Amended) A countermeasure method according to [any one of the preceding claims, characterised in that the choice of] claim 1, wherein the step of selecting the sequence to be executed is made at the start of execution of the algorithm by drawing a random value [(RND1), the chosen sequence is being the one used in each of the groups concerned].

8. (Amended) A countermeasure method according to [any one of the preceding claims, characterised in that the different] claim 1, wherein said manipulating means are tables of constants.

A COUNTERMEASURE METHOD IN AN ELECTRONIC COMPONENT USING A SECRET KEY CRYPTOGRAPHY ALGORITHM

The present invention relates to a countermeasure method in an electronic component using a secret key cryptography algorithm. They are used in applications where access to services or data is strictly controlled. They have an architecture formed around a microprocessor and memories, including a program memory which contains the secret key.

These components are notably used in chip cards, for certain applications thereof. These are for example applications involving access to certain data banks, banking applications, remote payment applications, for example for television, petrol dispensing or passing through motorway tolls.

These components or cards therefore use a secret key cryptography algorithm, the best known of which is the DES (standing for Data Encryption Standard in the British and American literature) algorithm. Other

secret key algorithms exist, such as the RC5 algorithm or the COMP128 algorithm. This list is of course not exhaustive.

In general terms and briefly, the function of these algorithms is to calculate an encoded message from a message applied as an input (to the card) by a host system (server, banking dispenser etc) and the secret key contained in the card, and to supply this encoded message in return to the host system, which for example enables the host system to authenticate the component or card, to exchange data, etc.

However, it has become clear that these components or cards are vulnerable to attacks consisting of a differential analysis of the current consumption and which enable ill-intentioned third parties to find the secret key. These attacks are referred to as DPA attacks, the English acronym for Differential Power Analysis.

The principle of these DPA attacks is based on the fact that the current consumption of the microprocessor executing the instructions varies according to the data being manipulated. Notably, an instruction from the microprocessor manipulating a data bit generates two different current profiles depending on whether this bit is "1" or "0". Typically, if the instruction is manipulating a "0", there is at this time of execution a first amplitude of the current consumed and if the instruction is manipulating a "1", there is a second amplitude of the consumed current, different from the first.

The characteristics of the cryptography algorithms are known: the calculations made, the parameters used. The only unknown is the secret key contained in the program memory. This cannot be derived solely from knowledge of the message applied as an input and the encoded message supplied in return.

However, in a cryptography algorithm, some calculated data depend only on the message applied in clear to the input of the card and the secret key contained in the card. Other data calculated in the algorithm can also be recalculated solely from the encoded message (generally supplied in clear at the output of the card to the host system) and the secret key contained in the card. More precisely, each bit of these particular data can be determined from the input or output message, and a limited number of particular bits of the key.

Thus, to each bit of a particular data item, there corresponds a sub-key formed by a particular group of bits of the key.

The bits of these particular data which can be predicted are hereinafter referred to as target bits.

The basic idea of the DPA attack is thus to use the difference in current consumption profile of an instruction depending on whether it is manipulating a "1" or a "0" and the possibility of calculating a target bit by means of the instructions of the algorithm using a known input or output message and a hypothesis on the corresponding sub-key.

The principle of the DPA attack is therefore to test a given sub-key hypothesis, applying, to a large number of current measurement curves, each relating to a known input message of the attacker, a Boolean selection function, a function of the sub-key hypothesis, and defined for each curve by the value predicted for a target bit.

By making an assumption on the sub-key concerned, it is in fact possible to predict the value "0" or "1" which this target bit will take for a given input or output message.

It is then possible to apply, as a Boolean selection function, the value, "0" or "1", predicted by the target bit for the sub-key hypothesis in question, in order to sort these curves into two packets: a first packet contains the curves which have seen the manipulation of the target bit at "0" and a second packet contains the curves which have seen the manipulation of the target bit at "1" according to the sub-key hypothesis. By taking the mean of the current consumption in each packet, a mean consumption curve M0(t) is obtained for the first packet and a mean consumption curve M1(t) for the second packet.

If the sub-key hypothesis is correct, the first packet actually contains all the curves amongst the N curves which have seen the manipulation of the target bit at "0" and the second packet actually contains all the curves amongst the N curves which have seen the manipulation of the target bit at "1". The mean consumption curve M0(t) of the first packet will then

have a mean consumption everywhere except at the times of execution of the critical instructions, with a current consumption profile characteristic of the manipulation of the target bit at "0" ($profile_0$). In other words, for all these curves, all the manipulated bits have had as many chances of equalling "0" as of equalling "1", except the target bit, which has always had the value "0". Which can be written:

$$M0(t) = [profile_0 + profile_1)/2]_{t \neq tci} + [profile_0]_{tci}$$

that is to say

$$M0(t) = [Vm_t]_{t \neq tci} + [profile_0]tci$$

where tci represents the critical instants, at which a critical instruction has been executed.

Likewise, the mean consumption curve M1(t) of the second packet corresponds to a mean consumption everywhere except at the times of execution of the critical instructions, with a current consumption profile characteristic of the manipulation of the target bit at "1" ($profile_1$). It is possible to write:

$$M1(t) = [profile_0 + profile_1)/2]_{t \neq tci} + [profile_1]_{tci}$$

that is to say

$$M1(t) = [Vm_t]_{t \neq tci} + [profile_1]tci$$

It has been seen that the two profiles, $profile_0$ and $profile_1$, are not equal. The difference between the curves M0(t) and M0(1) then gives a signal DPA(t),

whose amplitude is equal to $profile_0$-$profile_1$ at the critical instants tci of execution of the critical instructions manipulating this bit, that is to say, in the example depicted in Figure 1, at the places tc0 to tc6, and whose amplitude is approximately equal to zero outside the critical instants.

If the sub-key hypothesis is false, the sorting does not correspond to reality. Statistically, there is then in each packet as many curves which have actually seen the manipulation of the target bit at "0" as there are curves which have seen the manipulation of the target bit at "1". The resulting mean curve M0(t) is then situated around a mean value given by $(profile_0+profile_1)/2=Vm$, since, for each of the curves, all the bits manipulated, including the target bit, have as many chances of equalling "0" as of equalling "1".

The same reasoning on the second packet leads to a mean current consumption curve M1(t) whose amplitude is situated around a mean value given by $(profile_0+profile_1)/2=Vm$.

The signal DP(t) supplied by the difference M0(t)-M1(t) is in this case substantially equal to zero. The signal DPA(t) in the case of a false sub-key hypothesis is shown in Figure 2.

Thus the DPA attack exploits the difference in the current consumption profile during the execution of an instruction depending on the value of the bit manipulated, in order to effect a sorting of current consumption curves according to a Boolean selection

function for a given sub-key hypothesis. By effecting a differential analysis of the mean current consumption between the two packets of curves obtained, an information signal DPA(t) is obtained.

A DPA attack then consists overall in:

a- drawing N random messages (for example N equal to 1000);

b- having the algorithm executed by the card for each of the N random messages, reading the current consumption curve each time (measured on the supply terminal of the component);

c- making an assumption on a sub-key;

d- predicting, for each of the random messages, the value taken by one of the target bits whose value depends only the bits of the message (input or output) and on the sub-key taken as a hypothesis, in order to obtain the Boolean selection function;

e- sorting the curves according to this Boolean selection function (that is to say according to the value "0" or "1" predicted for this target bit for each curve under the sub-key hypothesis);

f- calculating, in each packet, the resulting mean current consumption curve;

g- taking the difference between these mean curves, in order to obtain the signal DPA(t).

If the hypothesis on the sub-key is correct, the Boolean selection function is correct and the curves of the first packet actually correspond to the curves for which the message supplied as an input or output gave a target bit at "0" in the card and the curves in the

second packet actually correspond to the curves for which the message applied as an input or output gave a target bit at "1" in the card.

Take the case in Figure 1: the signal DPA(t) is therefore not zero at times tc0 to tc6 corresponding to the execution of the critical instructions (those which manipulate the target bit). It suffices for there to have been at least one critical instant in the period of acquisition.

It should be noted that the attacker does not need to know precisely the critical instants.

If the sub-key hypothesis is not correct, the sorting does not correspond to reality and there are then in each packet as many curves corresponding in reality to a target bit at "0" as there are curves corresponding to a target bit at "1". The signal DPA(t) is substantially zero everywhere (the case shown in Figure 2). It is necessary to return to step c- and to make a new assumption on the sub-key.

If the hypothesis proves correct, it is possible to pass to the evaluation of other sub-keys, until the key has been reconstituted to the maximum possible extent. For example, with a DES algorithm, a key of 64 bits is used, of which only 56 are useful bits. With a DPA attack, it is possible to reconstitute at least 48 bits of the 56 useful bits.

The purpose of the present invention is to use, in an electronic component, a countermeasure method which gives rise to a zero signal DPA(t), even where the sub-key hypothesis is correct.

In this way, nothing makes it possible to distinguish the case of the correct sub-key hypothesis from the false sub-key hypotheses. By means of this countermeasure, the electronic component is protected against DPA attacks.

According to the invention, the countermeasure method makes the target bits, that is to say the data manipulated by critical instructions, unpredictable.

This is because, because of the countermeasure, for each message applied as an input, a target bit manipulated by a critical instruction takes the value 0 or 1 with equal probability. In each packet of curves which the attacker will make under a given sub-key hypothesis, by means of the Boolean selection function which he will have calculated, there will be as many curves actually having manipulated a target bit "0" as there are curves actually having manipulated a target bit at "1". The signal DPA(t) will always be zero, whether the sub-key hypothesis is correct or not.

In the invention, the concern is more particularly with the DES cryptography algorithm.

Such an algorithm comprises sixteen identical calculation rounds.

In such an algorithm, it has been possible to show that the data which can be predicted by an attacker are situated at the first round and at the last round, and that the critical instructions, in the sense of the DPA attack, are situated in the first three rounds and the last three rounds.

In the invention, a means of making unpredictable the data manipulated by these critical instructions in the first three and last three rounds, whilst obtaining the correct encoded message as an output, has in particular been sought.

One aim of the invention is therefore to make the data manipulated by the critical instructions unpredictable, whilst obtaining the correct final result (encoded message C).

One solution to these different technical problems has been found in the formation of a group (G1) comprising at least the first three rounds and another group (G4) comprising at least the last three rounds, and in the use of these groups of means for making unpredictable the data manipulated by the critical instructions contained in these rounds.

According to the invention, the results output from each group are correct.

As characterised, the invention therefore relates to a countermeasure method in an electronic component using a secret key cryptographic algorithm for calculating an encoded message from an input message, the use of the algorithm comprising sixteen calculation rounds, each round using first means for supplying an output data item from an input data item, the output data item and/or the derived data being manipulated by critical instructions in the first three and last three rounds. According to the invention, a group is formed comprising at least the first three rounds and another group comprising at least the last three rounds, and

with each of these groups there is associated a first sequence using the first means in each round and a second sequence using other means in certain rounds at least, the said first and second sequences being such that they supply the same result at the output from the last round in each group for the same given input message, the choice of the sequence to be executed in the groups concerned being a function of a statistical half probability distribution, in order to make all the data manipulated by the said critical instructions unpredictable.

In one embodiment, four groups each of four consecutive rounds are formed.

In another embodiment, two groups are formed, comprising respectively the first three and the last three rounds.

Other characteristics and advantages of the invention are detailed in the following description given for information and in no way limitatively, and with reference to the accompanying drawings, in which:

- Figures 1 and 2, already described, depict the signal DPA(t) which can be obtained as a function of a hypothesis on a sub-key of the secret key K, according to a DPA attack;

- Figures 3 and 4 are flow diagrams depicting the first rounds and the last rounds of the DES algorithm;

- Figure 5 is a block diagram of the operation SBOX used in the DES algorithm;

- Figure 6 shows an example of an elementary table of constants with one input and one output used in the operation SBOX;

- Figure 7 shows a first example of a flow diagram for executing the DES with a countermeasure method according to the invention;

- Figure 8 is a flow diagram of the first rounds of the DES according to a second sequence of the countermeasure method according to the first example depicted in Figure 7;

- Figures 9 and 10 depict respectively second and third elementary tables of constants used in the invention;

- Figure 11 depicts a second example of a flow diagram for executing the DES with a countermeasure method according to the invention;

- Figures 12 and 13 are flow diagrams of the first DES rounds respectively according to the second sequence and the first sequence of the countermeasure method according to the second example depicted in Figure 11;

- Figures 14 and 15 are flow diagrams relating to a third mode of applying the countermeasure method according to the invention;

- Figure 16 depicts a third elementary table of constants used in this fourth mode of application of the invention;

- Figure 17 depicts a flow diagram for execution of the DES according to a variant of the third mode of

application of the countermeasure method according to the invention; and

- Figure 18 depicts a simplified block diagram of a chip card containing an electronic component in which the countermeasure method according to the invention is implemented.

The DES secret key cryptographic algorithm (hereinafter the term DES or DES algorithm will more simply be used) contains 16 calculation rounds, denoted T1 to T16, as depicted in Figures 3 and 4.

The DES begins with an initial permutation IP on the input message M (Figure 3). The input message M is a word f of 64 bits. After permutation, a word e of 64 bits is obtained, which is divided into two in order to form the input parameters L0 and R0 of the first round (T1). L0 is a word d of 32 bits containing the 32 most significant bits of the word e. R0 is a word h of 32 bits containing the 32 least significant bits of the word e.

The secret key K, which is a word q of 64 bits, itself undergoes a permutation and compression in order to supply a word r of 56 bits.

The first round comprises an operation EXP PERM on the parameter R0, consisting of an expansion and a permutation, in order to supply a word l of 48 bits as an output.

This word l is combined with a parameter K1, in an operation of the exclusive OR type denoted XOR, in order to supply a word b of 48 bits. The parameter K1, which is a word m of 48 bits, is obtained from the word

r by a shift by one position (the operation denoted SHIFT in Figures 3 and 4) followed by a permutation and a compression (the operation denoted COMP PERM).

The word b is applied to an operation denoted SBOX, at the output of which a word a of 32 bits is obtained. This particular operation will be explained in more detail in relation to Figures 5 and 6.

The word a undergoes a permutation P PERM, giving as an output the word c of 32 bits.

This word c is combined with the input parameter L0 of the first round T1, in a logic operation of the exclusive OR type, denoted XOR, which supplies the word g of 32 bits as an output.

The word h (=R0) of the first round supplies the input parameter L1 of the following round (T2) and the word g of the first round supplies the input parameter R1 of the following round. The word p of the first round supplies the input r of the following round.

Other rounds T1 to T16 take place in a similar fashion, except with regard to the shift operation SHIFT, which takes place on one or two positions according to the rounds in question.

Each round Ti thus receives as an input the parameters Li-1, Ri-1 and r and supplies as an output the parameters Li and Ri and r for the following round Ti+1.

At the end of the DES algorithm (Figure 4), the encoded message is calculated from the parameters L16 and R16 supplied by the last round T16.

This calculation of the encoded message C comprises in practice the following operations:

- formation of a word e' of 64 bits by reversing the position of the words L16 and R16, and then concatenating them;

- application of the permutation IP^{-1} which is the reverse of that of the start of the DES, in order to obtain the word f' of 64 bits forming the encoded message C.

The operation SBOX is detailed in Figures 5 and 6. It comprises a table of constants TC_0 for supplying an output data item a as a function of an input data item b.

In practice, this table of constants TC_0 is in the form of eight elementary tables of constants $TC_0 1$ to $TC_0 8$, each receiving as an input only 6 bits of the word b, in order to supply only 4 bits of the word a as an output.

Thus the elementary table of constants $TC_0 1$ depicted in Figure 6 receives, as an input data item, the bits b1 to b6 of the word b and supplies as an output data item the bits a1 to a4 of the word a.

In practice these eight elementary tables of constants $TC_0 1$ to $TC_0 8$ are stored in the program memory of the electronic component.

In the operation SBOX of the first round T1, a particular bit of the output data item a of the table of constants TC_0 depends on only 6 bits of the data item b applied as an input, that is to say only 6 bits of the secret key K and the input message (M).

In the operation SBOX of the last round T16, a particular bit of the data item a output from the table of constants TC_0 can be recalculated from only six bits of the secret key K and the encoded message (C).

However, if the principle of the DPA attack is repeated, if a bit of the output data item a is chosen as the target bit, it suffices to make an assumption on 6 bits of the key K, in order to predict the value of a target bit for a given input (M) or output (C) message. In other words, for the DES, it suffices to make an assumption on a sub-key of 6 bits.

In a DPA attack on such an algorithm for a given target bit, it is therefore necessary to distinguish one correct sub-key hypothesis amongst 64 possible ones.

Thus, by taking only eight bits of the word a as target bits (one output bit per elementary table of constants $TC_0 1$ to $TC_0 8$), it is possible to discover up to 6x8=48 bits of the secret key, by making DPA attacks on each of these target bits.

In the DES, critical instructions in the sense of DPA attacks are therefore found at the start of the algorithm and at the end.

At the start of the DES algorithm, the data which can be predicted from an input message M and from a sub-key hypothesis are the data a and g calculated in the first round (T1).

The data item a from the first round T1 (Figure 3) is the output data item from the operation SBOX of the round in question. The data item g is calculated from

the data item a, by permutation (P PERM) and exclusive OR operation with the input parameter L0.

In fact, the data item c of the first round is a data item derived from the data item a of the first round. The derived data item c corresponds to a simple permutation of bits of the data item a.

The data item l of the first round is a data item derived from the data item g of the first round, since it corresponds to a permutation of the bits of the word g, certain bits of the word g also being duplicated.

Knowing a and g, it is also possible to know these derived data.

The critical instructions of the start of the algorithm are the critical instructions which manipulate either the data item which can be predicted, such as the data item a of the first round, or a derived data item.

The critical instructions manipulating the data item a of the first round T1 or the derived data item c are thus the instructions for the end of the operation SBOX, of the operation P PERM and the start of the operation XOR of the first round T1.

The critical instructions manipulating the data item g or the derived data are all the instructions of the end of the operation XOR of the end of the first round T1 as far as the instructions for the start of the operation SBOX of the second round T2, and the start of the XOR operation at the end of the third round T3 (L2 = h(T2) = g(T1)).

At the end of the DES algorithm, the data which can be predicted from an encoded message C and a sub-key hypothesis are the data item a of the sixteenth round T16 and the data item L15 equal to the word h of the fourteenth round T14.

The critical instructions manipulating the data item a of the sixteenth round or derived data are the instructions of the sixteenth round of the end of the operation SBOX, of the permutation operation P PERM and of the start of the operation XOR.

For the data item L15, the critical instructions manipulating this data item or derived data are all the instructions from the instructions of the end of the operation XOR of the fourteenth round T14, up to the instructions for the start of the operation SBOX of the fifteenth round T15, plus the instructions for the start of the operation XOR of the sixteenth round T16.

The countermeasure method according to the invention applied to this DES algorithm consists of having, for each critical instruction, as many chances for the critical instruction to manipulate a data item as its complement. Thus, whatever the target bit on which the DPA attack can be made, there are as many chances for the critical instructions manipulating this bit to manipulate a "1" or a "0".

In practice, this must be true for each of the potential target bits: in other words, the attacker having a choice between several possible attacks, that is to say between several possible Boolean selection functions for effecting a sorting of curves, for a

given sub-key hypothesis, the implementation of the countermeasure method according to the invention must ensure that the data manipulated by each of the critical instructions randomly take, half of the time, a value or its complement. With regard to the application of the countermeasure method according to the invention to the DES algorithm, it is therefore necessary to apply the countermeasure to the critical instructions for the start of the DES and to the critical instructions for the end of the DES, in order to be completely protected.

In the DES, all the data manipulated by critical instructions are an output item or data derived from an output data item from an operation SBOX.

This is because, at the start of the DES, the data which can be predicted are the data a and g of the first round T1. The data item a is the output data item of the operation SBOX of the first round. The data item g is calculated from the data item a, since g = P PERM(a) XOR L0. g is therefore a data item derived from the output data item a of the operation SBOX of the first round. Thus all the data manipulated by the critical instructions of the start of the DES result directly or indirectly from the output data item a of the operation SBOX of the first round.

With regard to the end of DES, the data which can be predicted are the data item a of the sixteenth round T16 and the data item g of the fourteenth round T14, g being equal to L15.

The data item a is the output data item of the operation SBOX of the sixteenth round T16.

As for the data item L15, this is calculated, in the normal execution of the DES algorithm, from the output data item a of the operation SBOX of the fourteenth round T14: L15 = P PERM(a) XOR L14.

If the output data a of these particular operations SBOX are made unpredictable, all the derived data are also made unpredictable: all the data manipulated by the critical instructions of the DES algorithm are therefore made unpredictable. If it is considered that these operations SBOX constitute first means for supplying an output data item S=a from an input data item E=b, the countermeasure method applied to the DES algorithm consists of using other means for making the output data item unpredictable, so that this output data item and/or derived data manipulated by the critical instructions are all unpredictable.

According to the invention, a group formed by at least the first three rounds and another group formed by at least the last three rounds are formed. These groups therefore contain all the rounds comprising the critical instructions.

With these two groups there is associated a first sequence using the first means for all the rounds and a second sequence using the other means for at least some rounds.

In the other rounds which are not in these groups, it is possible to continue to use the first means.

The use of these other means is such that the output result, that is to say the encoded message, remains correct.

These other means can comprise several different means. They are such that they make the complemented data item correspond to one or other or both data items amongst the input and output data of the first means.

Thus, considering a large execution number, the groups will use the first sequence, which is the normal sequence of the algorithm, on average half of the time, and the other sequence half of the time. The data manipulated by the critical instructions in these groups, corresponding to certain intermediate results, will therefore on average be complemented half of the time. On a large number of curves there will therefore be statistically as many chances that a given target bit will be at 1 or at 0.

Figure 7 depicts a first embodiment of the invention.

In this embodiment, the sixteen rounds of the DES algorithm are divided into four groups G1 to G4 of four successive rounds. The group G1 thus comprises the rounds T1 to T4, the group G2 the rounds T5 to T8, the group G3 the rounds T9 to T12 and the group G4 the rounds T13 to T16.

Two sequences are associated with each group. A first sequence SEQA consists of using the first means TC_0 for each round. A second sequence SEQB consists of using other means for at least some rounds.

In the example depicted, these other means comprise second means TC_2 and third means TC_1.

The second means TC_2 are used in the second round and the penultimate round of each group; that is to say in T2, T3 of G1, T6, T7 of G2, T10, T11 of G3 and T14 and T15 of G4.

The third means TC_1 are used in the first round and last round of each group. That is to say in T1, T4 of G1, T5, T8 of G2, T9, T12 of G3 and T13, T16 of G4.

In practice, these different means are tables of constants. The first means correspond to the first table of constants TC_0, corresponding to the normal execution of the DES. The other means TC_1 and TC_2 are defined with respect to this first table of constants TC_0, by complementation.

The second means TC_2 are such that, for the complement /E of the input data item E, they supply the complement of the output data item S of the first means TC_0. One example of a second elementary table TC_21 corresponding to the first elementary table of constants TC_01 is depicted in Figure 9. It should be noted that the notation of the complement /E used in the text, corresponds the notation with a bar above the complemented data item in the drawings.

The third means are such that, for the input data item E, they supply the complement /S of the output data item S of the first means TC_0. One example of a third elementary table TC_11 corresponding to the first elementary table of constants TC_01 is depicted in Figure 10.

The calculation program then consists, at the start of the execution of the algorithm, in drawing a random value RND1 equal to 0 or to 1, and then testing this value RND1. In the example, if RND1 is equal to 1, the calculation is made using the second sequence SEQB for each group G1 to G4.

If RND1 is equal to 0, the calculation is made using the first sequence SEQA for each group.

Whether the first or second sequence is used, the correct result for the output parameters is obtained at the output of each group. Thus the output parameters L4 and R4 of the first group G1, L8 and R8 of the second group G2, L12 and R12 of the third group G3, L16 and R16 of the fourth group G4 are correct whatever the sequence used.

When all the rounds have been executed, the correct parameters L16 and R16, which will make it possible to calculate the correct encoded message C, are obtained.

On the other hand, within the groups, certain intermediate results do not have the same values according to the sequence used, but complementary values, as will be shown with reference to Figures 3 and 8.

Figure 3, already described, corresponds in fact to the flow diagram for calculating the four rounds T1, T2, T3 and T4 of the first group G1, in the first sequence SEQA.

Figure 8 shows the detailed flow diagram of the four rounds T1, T2, T3 and T4 of the first group G1, in the second sequence SEQB.

In the second sequence, the round T1 uses the third means TC_1. At the output of the operation SBOX, the data item /a (Figure 8) is therefore obtained, instead of the data item a with the first sequence SEQA (Figure 3).

The operation P PERM of the round T1, which is a simple permutation, will therefore also supply as an output a complemented data item /c with respect to the sequence SEQA.

The data item g, which is obtained by an exclusive OR between a complemented data item /c and a non-complemented data item L0, will also supply as an output a complemented data item /g.

Thus, with the third means of the round T1, all the following complemented data are obtained, with respect to the data which would be obtained with the sequence SEQA:

- in the round T1: /a, /c, /g;
- in the round T2: /R1, /h, /l, /b;
- in the round T3: /L2.

The second means TC_2 used in the round T2 are then arrived at. According to their definition, by applying the complemented data item /b, the complemented data item /a is obtained as an output. By taking this reasoning as far as the end of the round T4, noting that an exclusive OR between two complemented data gives a non-complemented result (for example /L3 XOR /c

= g in the round T4), the non-complemented data L4, R4 are obtained at the output of the round T4.

In addition, it is found that, for all the critical instructions for the start of the DES, the critical instructions will manipulate, in a random manner according to the data item RND1, the data or their complements depending on whether the sequence executed is the first SEQA or the second SEQB.

The countermeasure method, in this first embodiment, is therefore of great interest. It requires only two additional operations in the DES calculation program, the drawing of the random value and the testing of this value. The program memory for its part must contain the three different means used, that is to say the three tables of constants TC_0, TC_1, TC_2.

Returning to Figure 7, it can be noted that there is no need for a countermeasure in the groups in the middle G2 and G3, since they do not contain any critical instructions within the meaning of a DPA attack. It would therefore be possible to apply the countermeasure method with its two sequences SEQA and SEQB only to the first and second groups G1 and G4. It would suffice to apply the first sequence SEQA systematically to groups G2 and G3.

However, applying the countermeasure method to all the groups gives consistency to the whole.

Thus the two sequences SEQA and SEQB are preferably associated with each of the groups G1 to G4.

A second embodiment of the countermeasure method according to the invention is depicted in Figure 11. This second embodiment is in fact a variant of the first. The advantage of this variant is using, as other means in the sequence SEQB, only the second means TC_2. This is because it has been seen that the different means TC_0, TC_1, TC_2 correspond in practice to tables of constants each comprising eight elementary tables of constants, which occupies a not insignificant amount of space in the program memory.

This variant therefore consists in using the two means TC_2 in the sequence SEQB. For this, an additional operation CP, for complementing the input data item applied to the second means, is provided in the program for calculating the first and second rounds of each group. This additional operation CP is in practice an exclusive OR of the input data item with logic ones. If reference is made to Figure 12 depicting the detailed flow diagram of the second sequence SEQB for calculating the four rounds T1 to T4 of the first group G1, it is a case of complementing the data item b before applying it to the input of the operation SBOX of the rounds T1 and T4. As the second means TC2 complement the input, the complementation operation CP plus the second means TC_2 are equivalent to the third means TC_1 used in the first embodiment of the invention, that is to say to a data item which is not complemented at the input.

However, for the countermeasure method according to the second embodiment to be effective, it is

necessary for the number of instructions to be exactly the same whatever the calculation sequence used. This is because, if any difference existed between the two possible sequences SEQA and SEQB, there would then be a possibility of a successful DPA attack.

For this reason, and as depicted in Figure 13, there is provided, in the rounds T1 and T4 of the first sequence SEQA, an operation ID of copying identically, which consists of an exclusive or with logic zeros at the input of the operation SBOX, in order not to modify the input data item whilst applying the same instructions as for the additional operation CP.

In this way, there are the same number of instructions in the two sequences.

Figure 14 depicts a third embodiment of the countermeasure method according to the invention.

In this embodiment, a first group G1 is formed with the first three rounds T1, T2, T3 and another group G4 with the last three rounds T14, T15, T16. There is associated with each group a first sequence SEQA using the first means TC_0 for each round and a second sequence using other means for at least some rounds.

At the output of each group G1, G4, the correct output result L3, R3 and L16, R16 is obtained, whatever the sequence SEQA or SEQB used.

The other means are, in the example, the third means TC_1 already seen in relation to the first embodiment and the fourth means TC_3.

These fourth means TC_3 are defined, with respect to the first means TC_0, as making the output data item S correspond to the complement /E of the input data item E. A corresponding elementary table of constants TC_31 is depicted in Figure 16.

For the other rounds not included in the groups, that is to say for the rounds T4 to T13, the first means TC_0 are applied.

Thus, after having drawn the random value RND1, this value is tested in order to determine the sequence to be applied to the first group, and the output continues with the parameters L3, R3 calculated, the following rounds being executed with the first means TC_0. At the end of the round T13, the sequence determined by the random value RND1 is applied to the group G4. The parameters L16, R16 are obtained, which will serve to calculate the encoded message C.

Figure 15 is a corresponding detailed flow diagram, for the second sequence SEQB.

It is clear from this flow diagram that complemented data are obtained (the complementation being denoted by a bar above the data item) for all the critical instructions of these rounds. And the data L3 and R3 at the output of the third round are not complemented. The execution of the algorithm can be continued by passing to the round T4, at which the first means TC_3 are applied according to the normal execution of the algorithm.

In this figure, it can be remarked that, in the operation SBOX of the third round T3 it would be

possible to use the first means TC_0 instead of the third means TC_1, providing an additional complementation operation CP at the output of the operation SBOX. This is an equivalent solution.

It is then necessary to make the additional operation of identical copying ID in the sequence SEQA correspond to this additional complementation operation in the sequence SEQB.

Figure 17 depicts an execution flow diagram using this variant. For the third round in the two groups G1 and G4, in the first sequence SEQA, the first means TC_0 are used, followed at the output by the additional copying operation ID, which is denoted $T3(TC_0, ID)$. In the second sequence SQB, for the third round the first means TC_0 are used, followed at the output by the additional complementation operation CP, which is denoted $T3(TC_0, CP)$.

Thus the second embodiment and this variant of the third embodiment show the use of additional operations at the input or output of the different means.

input or output of the means used. To each additional complementation operation CP in the second sequence there then corresponds an additional operation of identical copying ID in the first sequence SEQA.

The present invention applies to the DES secret key cryptography algorithm, for which several examples of non-limitative applications have been described. It applies more generally in a secret key cryptography algorithm with sixteen calculation rounds, the critical instructions of which are situated amongst the situations of the first three or last three rounds.

A electronic component 1 using a countermeasure method according to the invention in a DES secret key cryptography algorithm typically comprises, as shown in Figure 18, a microprocessor μP, a program memory 2 and a working memory 3. In order to be able to manage the use of the different means TC_0, TC_1, TC_2 according to the invention, which are, in practice, tables of constants stored in the program memory, means 4 of generating a random value between 0 and 1 are provided which, if reference is made to the flow diagrams in Figures 7 and 11, supply the value of RND1 at each execution of the DES. Such a component can in particular be used in a chip card 5, for improving their resistance to tampering.

CLAIMS

1. A countermeasure method in an electronic component using a secret key cryptographic algorithm (K) for calculating an encoded message (C) from an input message (M), the use of the algorithm comprising sixteen calculation rounds (T1, …, T16), each round using first means (TC_0) for supplying an output data item from an input data item, the output data item and/or the derived data being manipulated by critical instructions in the first three (T1, T2, T3) and last three (T14, T15, T16) rounds, characterised in that a group (G1) is formed comprising at least the first three rounds and another group (G4) comprising at least the last three rounds, and in that with each of these groups (G1 and G4) there is associated a first sequence (SEQA) using the first means (TC_0) in each round and a second sequence (SEQB) using other means (TC_1, TC_2, TC_3) in certain rounds at least, the said first and second sequences being such that they supply the same result at the output from the last round in each group for the same given input message (M), the choice of the sequence to be executed in the groups concerned being a function of a statistical half probability distribution, in order to make all the data manipulated by the said critical instructions unpredictable.

2. A countermeasure method according to Claim 1, characterised in that the other means are such that they complement one or other or both of the input (E) and/or output (S) data of the first means.

3. A countermeasure method according to Claim 2, characterised in that the second sequence (SEQB) comprises, for one or more rounds, an additional complementation operation (CP) at the input or output of the means used, and in that, to each additional complementation operation in the second sequence there corresponds an additional operation of identical copying (ID) in the first sequence (SEQA).

4. A countermeasure method according to any one of the preceding claims, characterised in that four groups (G1, ... G4) of each of four successive rounds (T1, ... T4) are formed, in that the first sequence (SEQA) is associated with each group and in that the second sequence (SEQB) is associated at least with the first group (G1) and the last group (G4).

5. A countermeasure method according to Claim 4, characterised in that the second sequence (SEQB) is associated with each of the groups (G1, ... G4).

6. A countermeasure method according to any one of Claims 1 to 3, characterised in that the first group (G1) is formed by the first three rounds (T1, T2, T3) and in that the last group is formed by the last three rounds (T14, T15, T16).

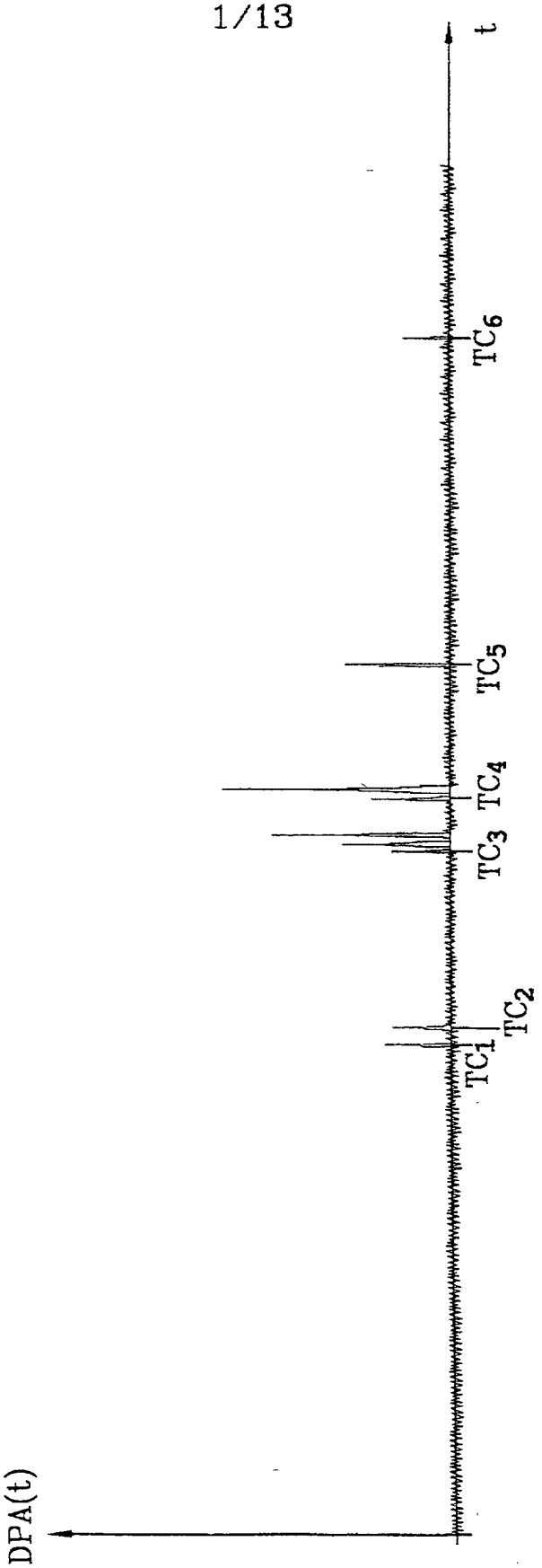
7. A countermeasure method according to any one of the preceding claims, characterised in that the choice of the sequence to be executed is made at the start of execution of the algorithm by drawing a random value (RND1), the chosen sequence being the one used in each of the groups concerned.

8. A countermeasure method according to any one of the preceding claims, characterised in that the different means are tables of constants.

9. An electronic security component using the countermeasure method according to any one of the preceding claims, characterised in that the different means (TC_0, TC_1, TC_2) for supplying an output data item from an input data item are fixed in the program memory of the said component and in that it comprises means (4) of generating a random value (RND1) at 0 or 1 in order to manage the use of the said different means.

10. A chip card comprising an electronic security component according to Claim 9.

FIG.1



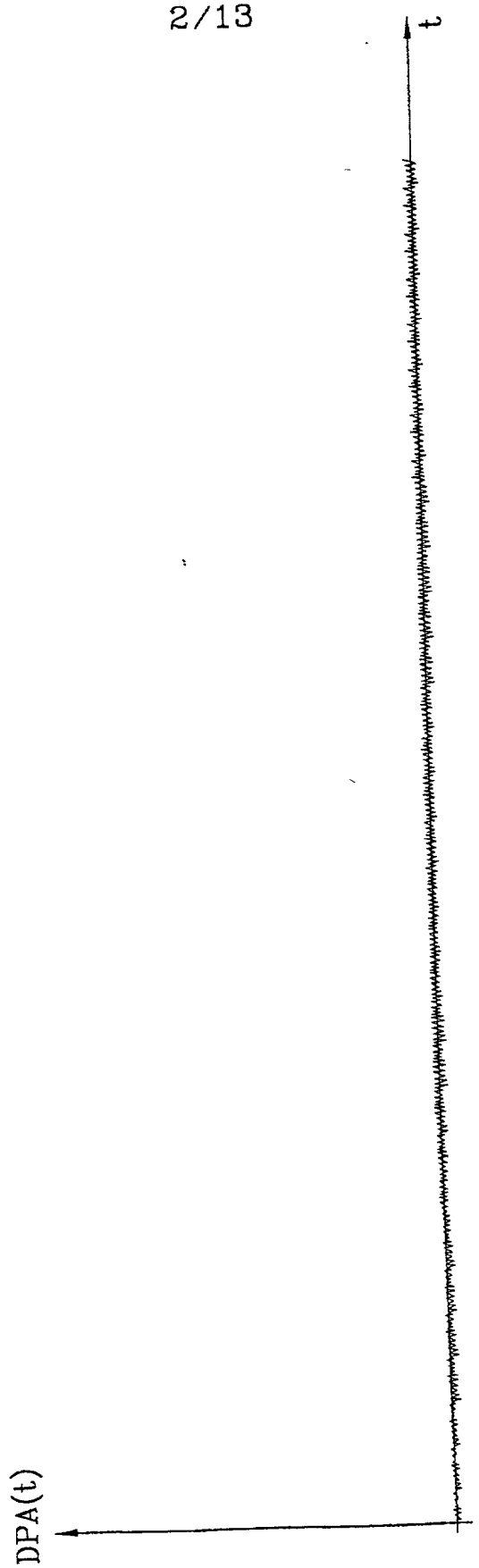


FIG.2

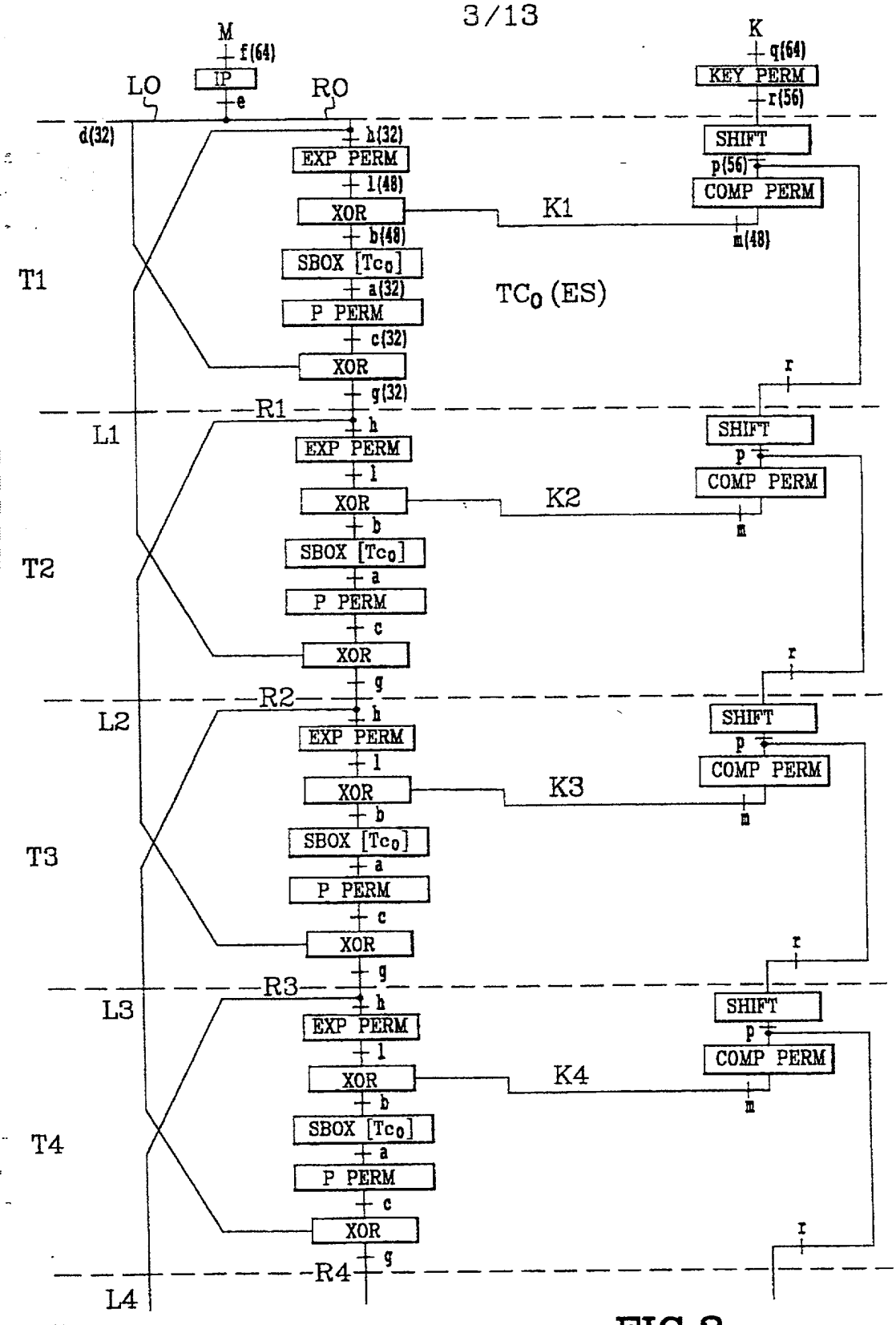

M
f(64)
IP
e
LO
RO
K
q(64)
KEY PERM
r(56)
d(32)
h(32)
EXP PERM
l(48)
XOR
K1
b(48)
SBOX [Tc0]
a(32)
P PERM
c(32)
XOR
g(32)
SHIFT
p(56)
COMP PERM
m(48)
TC0 (ES)
T1
R1
L1
T2
R2
L2
T3
R3
L3
T4
R4
L4
T5
K2
K3
K4

FIG.3

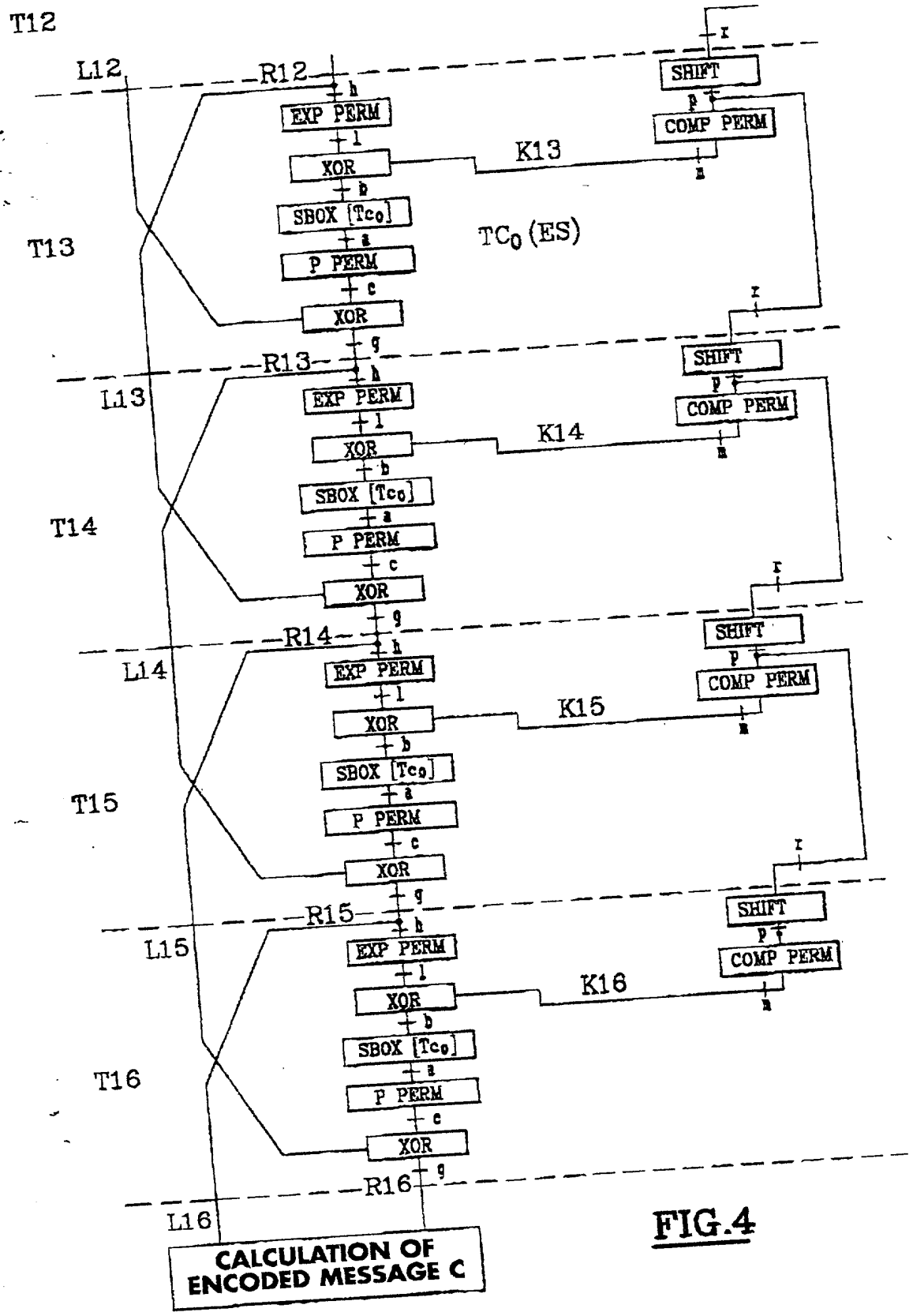

T12
L12
R12
EXP PERM
XOR
K13
SBOX [Tc0]
TC0 (ES)
T13
P PERM
XOR
SHIFT
COMP PERM
R13
L13
EXP PERM
XOR
K14
SBOX [Tc0]
T14
P PERM
XOR
SHIFT
COMP PERM
R14
L14
EXP PERM
XOR
K15
SBOX [Tc0]
T15
P PERM
XOR
SHIFT
COMP PERM
R15
L15
EXP PERM
XOR
K16
SBOX [Tc0]
T16
P PERM
XOR
SHIFT
COMP PERM
R16
L16
CALCULATION OF ENCODED MESSAGE C

FIG.4

FIG.5

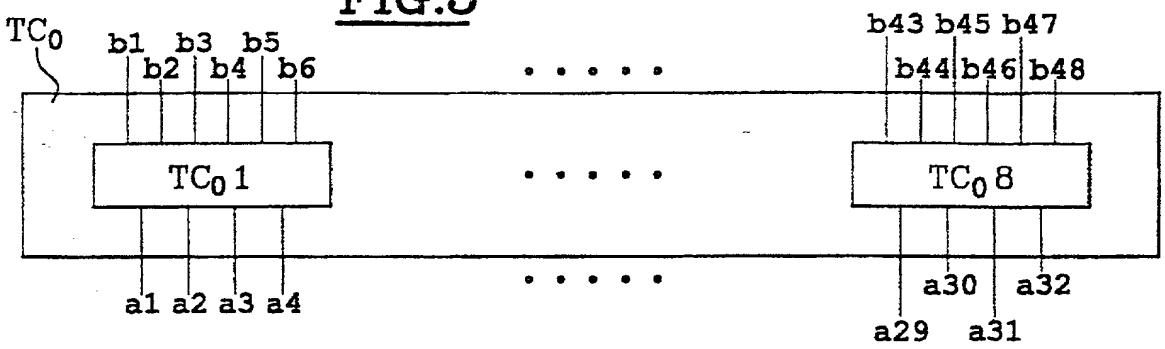


FIG.6

TC0 1

E=b1b2b3b4b5b6	S=a1a2a3a4
000000	1101
000001	0101
⋮	⋮
111111	1010

FIG.10

TC1 1

E=b1b2b3b4b5b6	/S=a1a2a3a4
000000	0010
000001	1010
⋮	⋮
111111	0101

FIG.9

TC2 1

/E=b1b2b3b4b5b6	/S=a1a2a3a4
000000	0101
⋮	⋮
111110	1010
111111	0010

FIG.16

TC3 1

/E=b1b2b3b4b5b6	S=a1a2a3a4
000000	1010
⋮	⋮
111110	0101
111111	1101

FIG.18

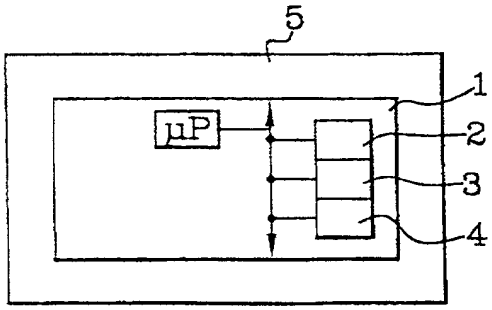
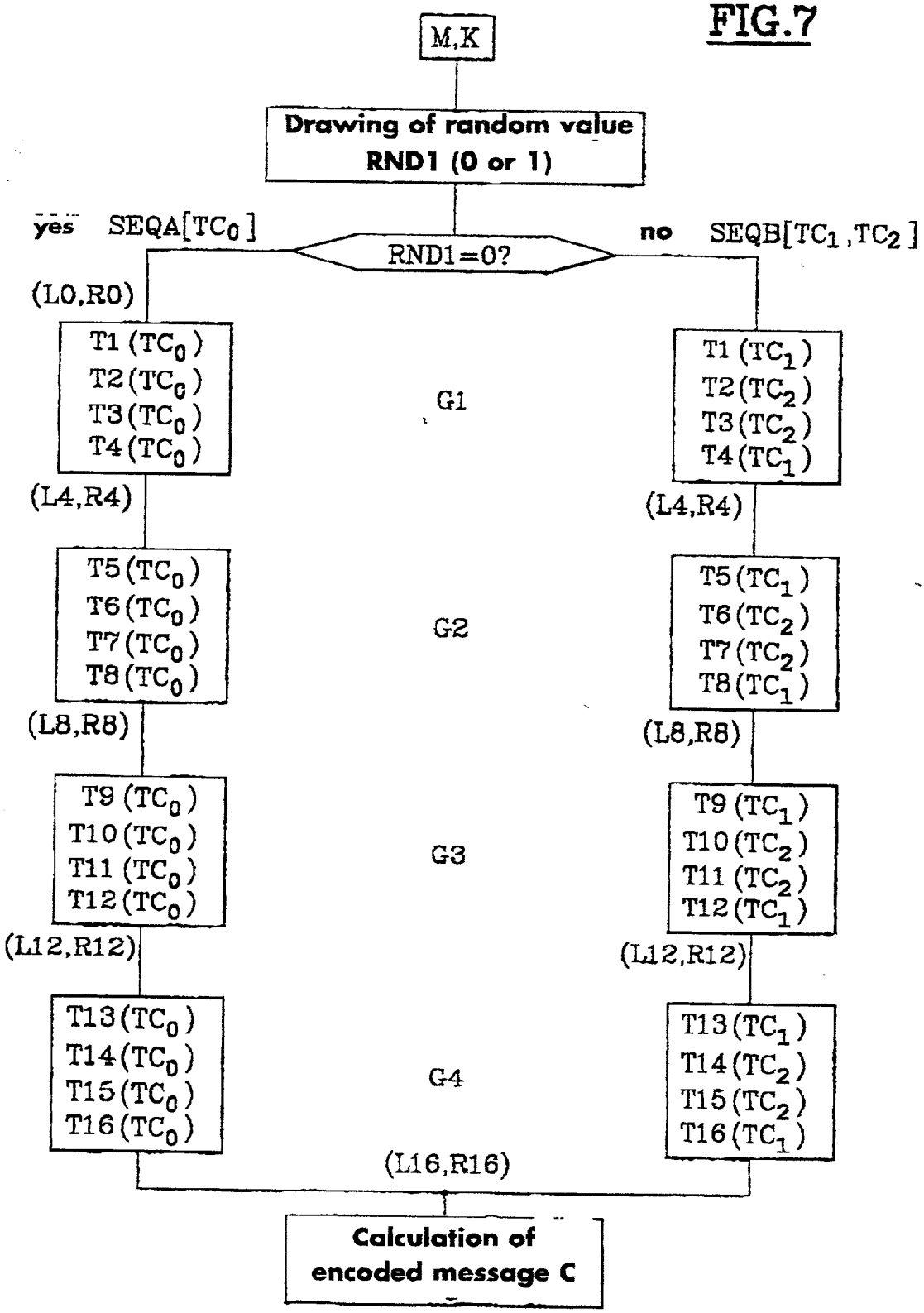


FIG.7


M,K
Drawing of random value RND1 (0 or 1)
RND1=0?
yes SEQA[TC0]
no SEQB[TC1,TC2]
(L0,R0)
T1(TC0) T2(TC0) T3(TC0) T4(TC0)
T1(TC1) T2(TC2) T3(TC2) T4(TC1)
G1
(L4,R4)
(L4,R4)
T5(TC0) T6(TC0) T7(TC0) T8(TC0)
T5(TC1) T6(TC2) T7(TC2) T8(TC1)
G2
(L8,R8)
(L8,R8)
T9(TC0) T10(TC0) T11(TC0) T12(TC0)
T9(TC1) T10(TC2) T11(TC2) T12(TC1)
G3
(L12,R12)
(L12,R12)
T13(TC0) T14(TC0) T15(TC0) T16(TC0)
T13(TC1) T14(TC2) T15(TC2) T16(TC1)
G4
(L16,R16)
Calculation of encoded message C

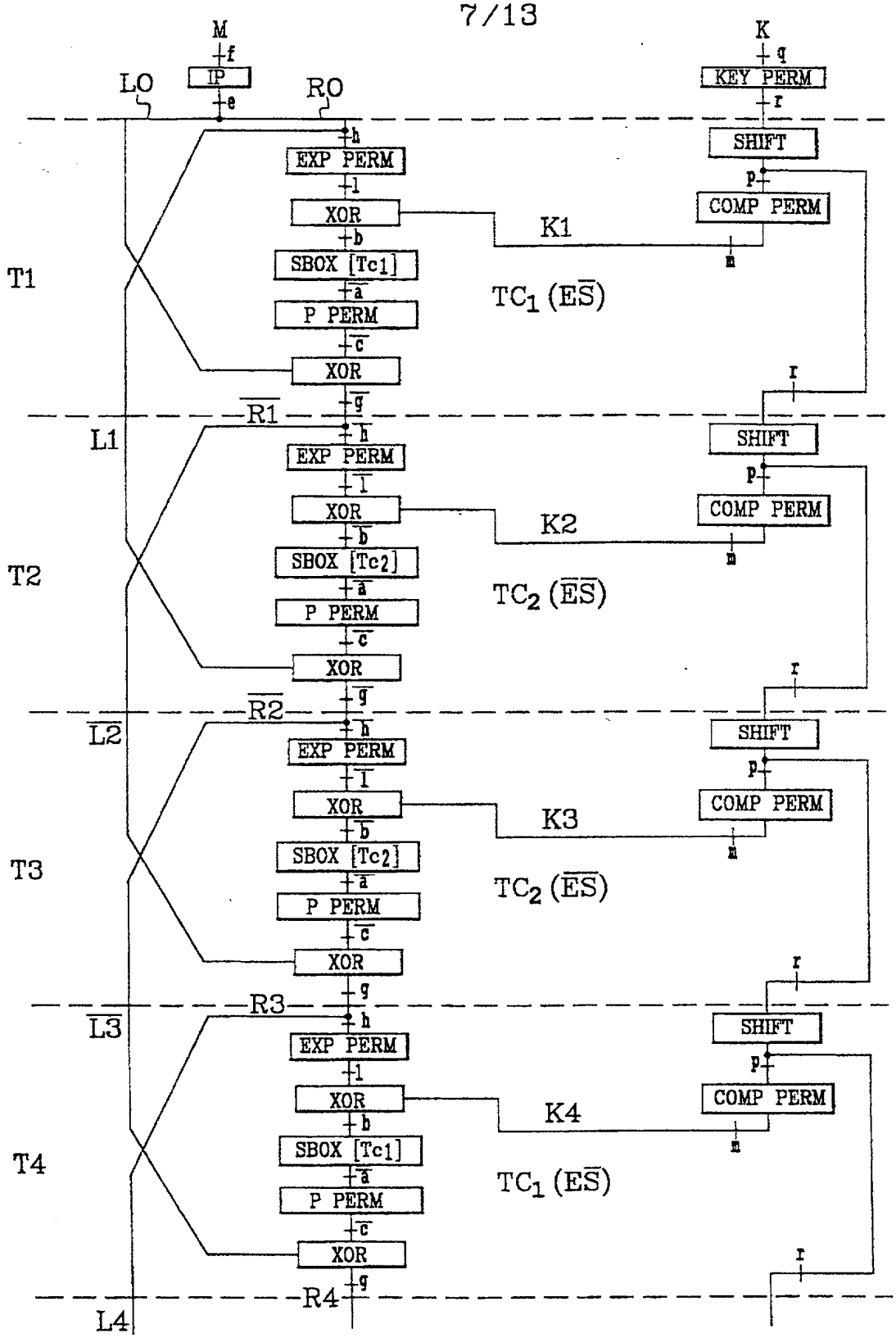
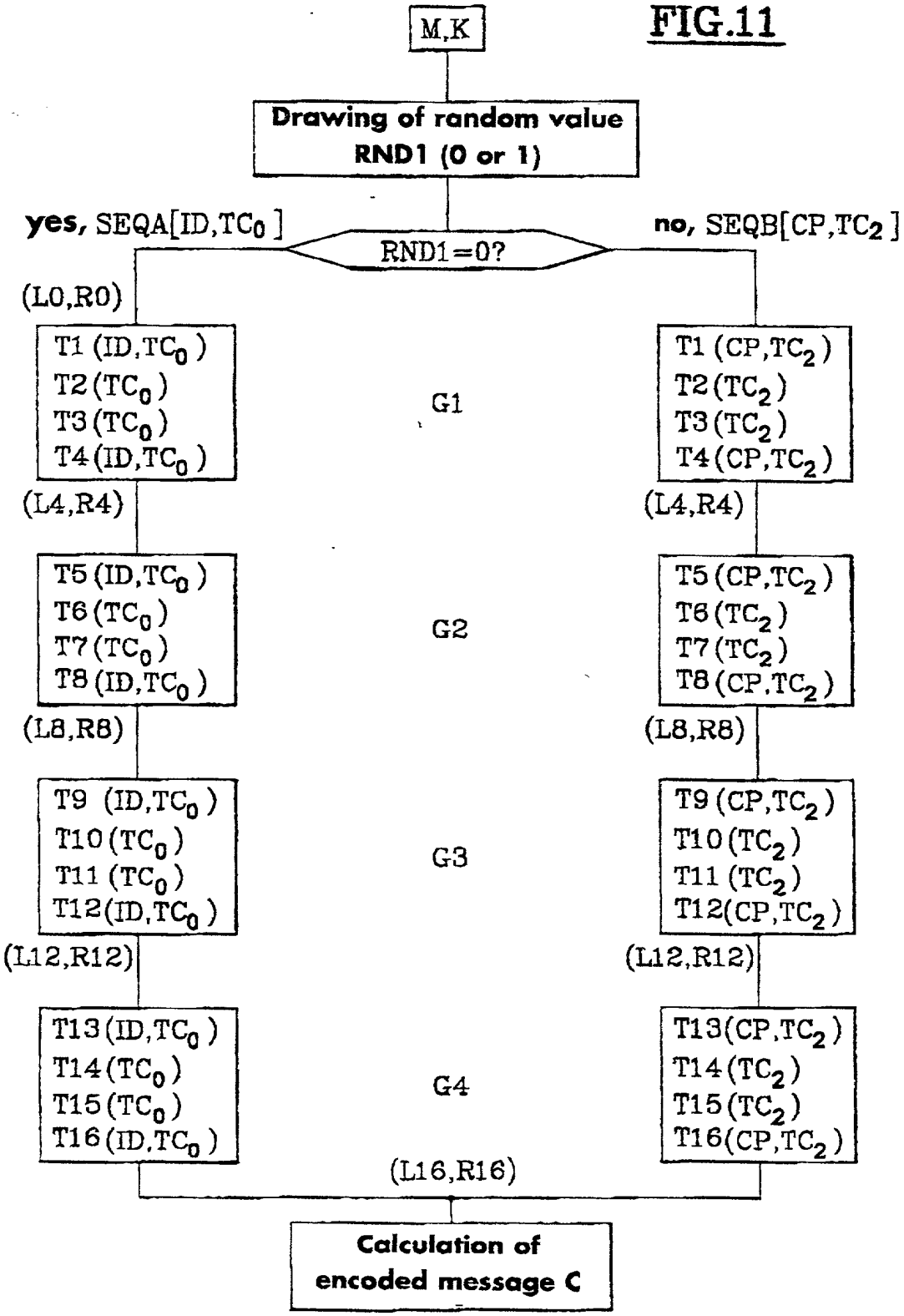

M
f
IP
e
L0
R0
K
q
KEY PERM
r
h
EXP PERM
l
XOR
b
SBOX [Tc1]
a
P PERM
c
XOR
g
SHIFT
p
COMP PERM
m
K1
T1
TC1 (ES)
R1
L1
K2
T2
TC2 (ES)
R2
L2
K3
T3
TC2 (ES)
R3
L3
K4
T4
TC1 (ES)
R4
L4

FIG.8

FIG.11


M,K
Drawing of random value RND1 (0 or 1)
yes, SEQA[ID,TC0]
RND1=0?
no, SEQB[CP,TC2]
(L0,R0)
T1 (ID,TC0)
T2(TC0)
T3(TC0)
T4(ID,TC0)
G1
T1 (CP,TC2)
T2(TC2)
T3(TC2)
T4(CP,TC2)
(L4,R4)
(L4,R4)
T5(ID,TC0)
T6(TC0)
T7(TC0)
T8(ID,TC0)
G2
T5(CP,TC2)
T6(TC2)
T7(TC2)
T8(CP,TC2)
(L8,R8)
(L8,R8)
T9 (ID,TC0)
T10(TC0)
T11 (TC0)
T12(ID,TC0)
G3
T9(CP,TC2)
T10(TC2)
T11 (TC2)
T12(CP,TC2)
(L12,R12)
(L12,R12)
T13(ID,TC0)
T14(TC0)
T15(TC0)
T16(ID,TC0)
G4
T13(CP,TC2)
T14(TC2)
T15(TC2)
T16(CP,TC2)
(L16,R16)
Calculation of encoded message C

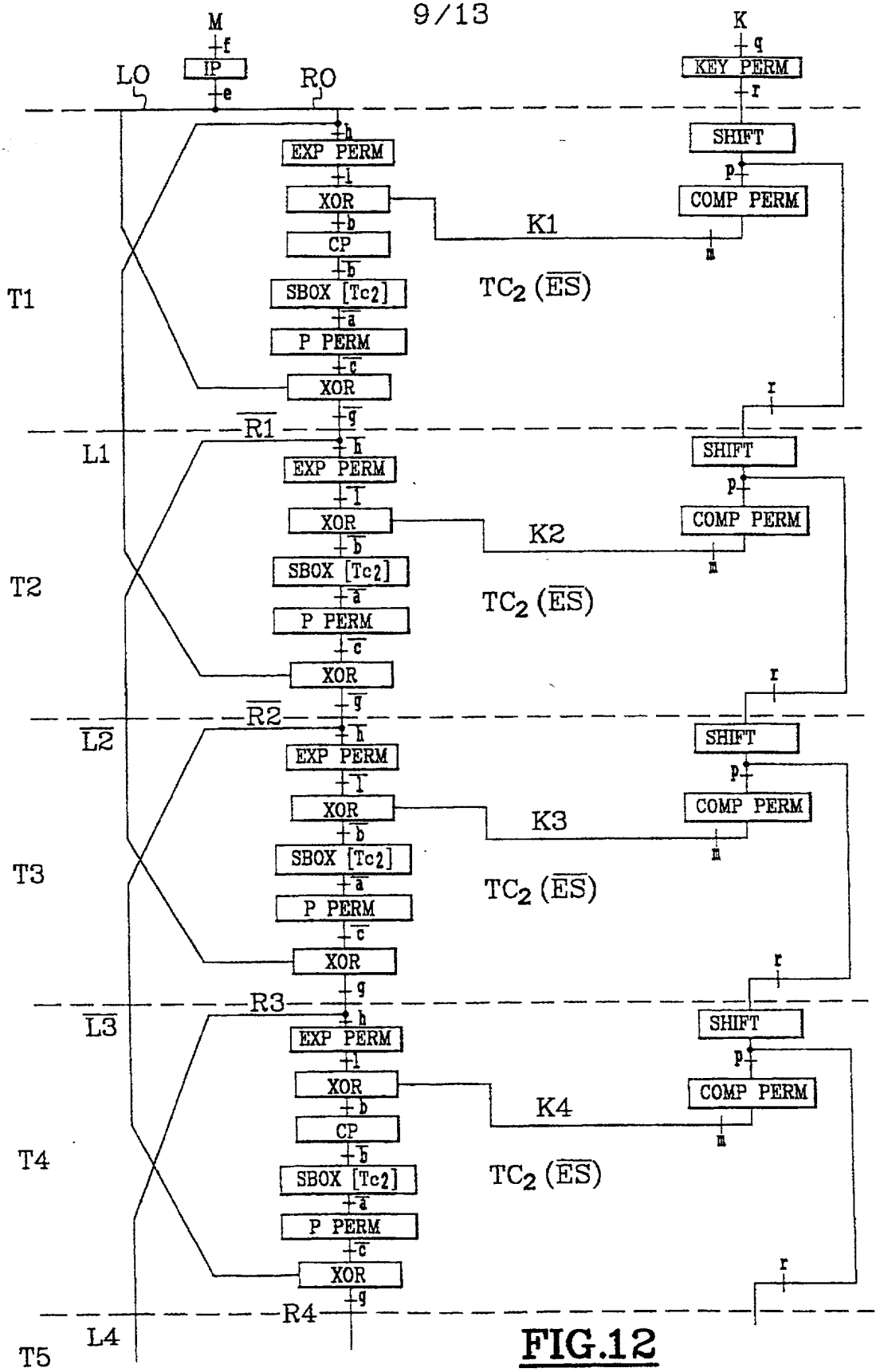

M
f
IP
e
LO
RO
K
q
KEY PERM
r
SHIFT
p
COMP PERM
m
K1
h
EXP PERM
l
XOR
b
CP
b
SBOX [Tc2]
a
P PERM
c
XOR
g
TC2 (ES)
T1
R1
L1
K2
T2
R2
L2
K3
T3
R3
L3
K4
T4
R4
L4
T5

FIG.12

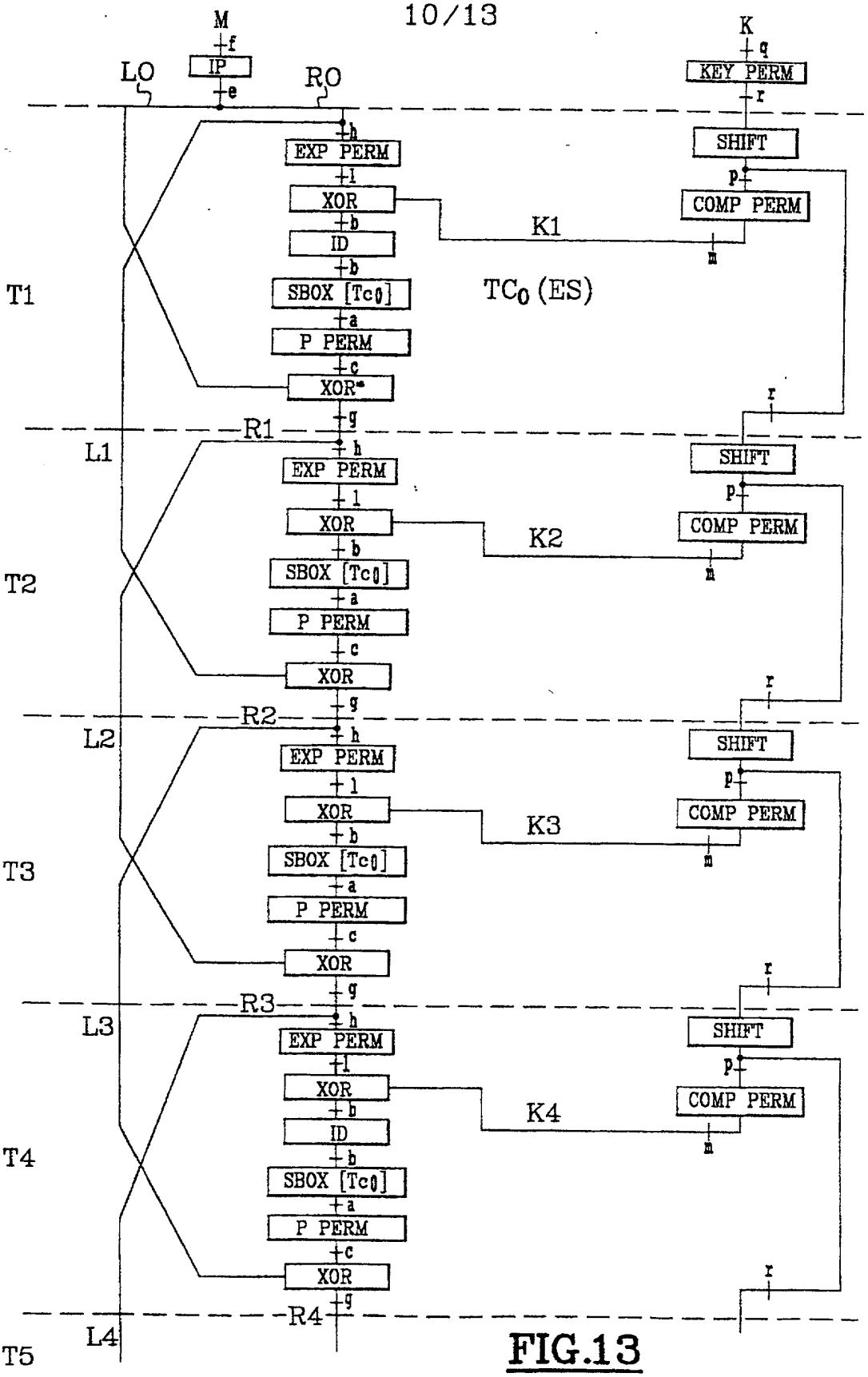

M
f
IP
e
L0
R0
K
q
KEY PERM
r
h
EXP PERM
SHIFT
l
XOR
p
COMP PERM
K1
b
ID
m
b
T1
SBOX [Tc0]
TC0 (ES)
a
P PERM
c
XOR
g
r
R1
L1
h
EXP PERM
SHIFT
l
XOR
p
COMP PERM
K2
b
m
T2
SBOX [Tc0]
a
P PERM
c
XOR
g
r
R2
L2
h
EXP PERM
SHIFT
l
XOR
p
COMP PERM
K3
b
m
T3
SBOX [Tc0]
a
P PERM
c
XOR
g
r
R3
L3
h
EXP PERM
SHIFT
l
XOR
p
COMP PERM
K4
b
ID
m
T4
b
SBOX [Tc0]
a
P PERM
c
XOR
g
r
R4
L4
T5

FIG.13

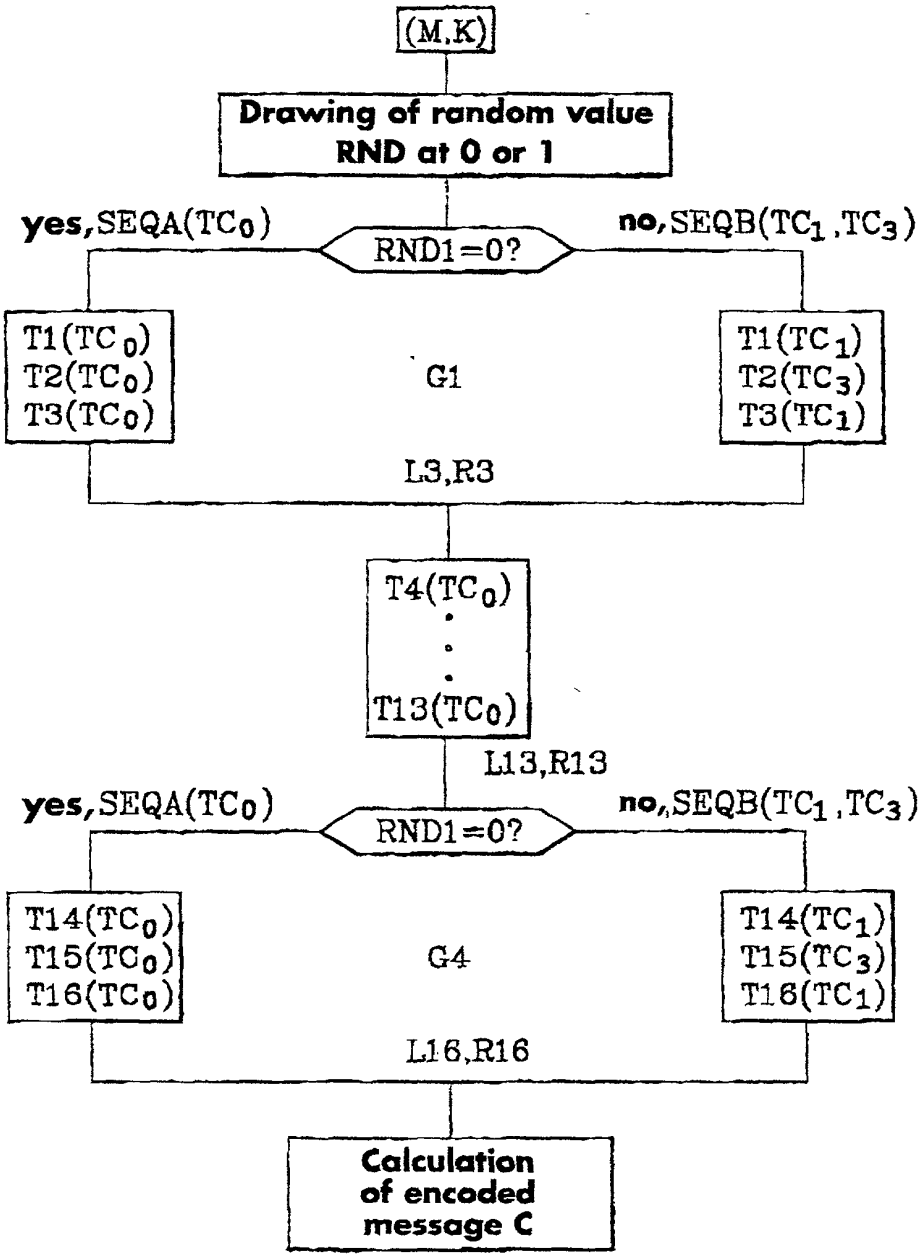

(M,K)
Drawing of random value RND at 0 or 1
yes, SEQA(TC0)
RND1=0?
no, SEQB(TC1,TC3)
T1(TC0) T2(TC0) T3(TC0)
G1
T1(TC1) T2(TC3) T3(TC1)
L3,R3
T4(TC0)
T13(TC0)
L13,R13
yes, SEQA(TC0)
RND1=0?
no, SEQB(TC1,TC3)
T14(TC0) T15(TC0) T16(TC0)
G4
T14(TC1) T15(TC3) T16(TC1)
L16,R16
Calculation of encoded message C

FIG.14

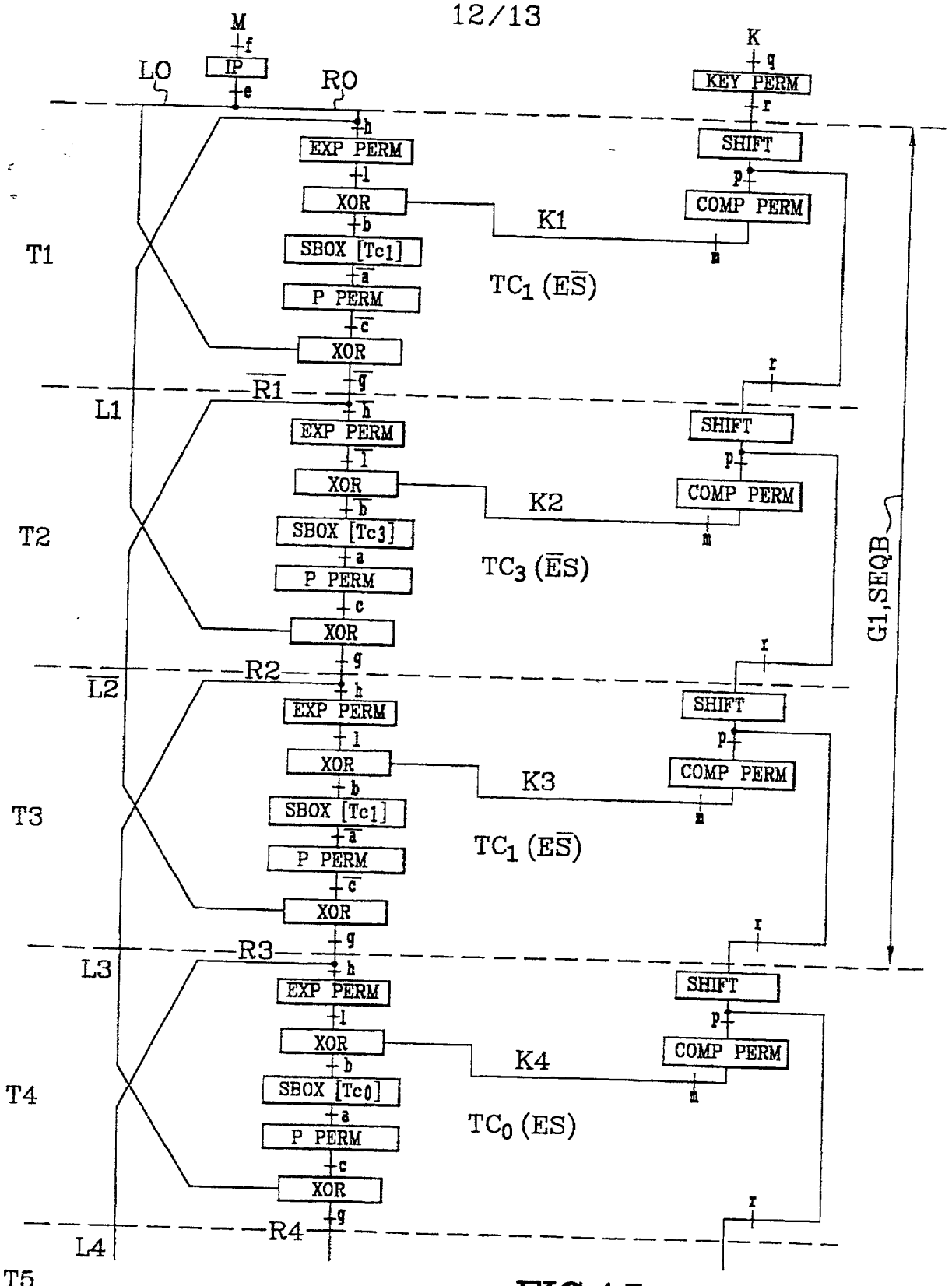

M
f
IP
e
L0
R0
K
q
KEY PERM
r
h
EXP PERM
SHIFT
l
XOR
p
COMP PERM
K1
b
SBOX [Tc1]
n
T1
a
P PERM
TC1 (ES)
c
XOR
g
R1
r
L1
h
EXP PERM
SHIFT
l
XOR
p
COMP PERM
K2
b
SBOX [Tc3]
m
T2
a
P PERM
TC3 (ES)
c
XOR
G1,SEQB
g
r
R2
L2
h
EXP PERM
SHIFT
l
XOR
p
COMP PERM
K3
b
SBOX [Tc1]
n
T3
a
P PERM
TC1 (ES)
c
XOR
g
r
R3
L3
h
EXP PERM
SHIFT
l
XOR
p
COMP PERM
K4
b
SBOX [Tc0]
m
T4
a
P PERM
TC0 (ES)
c
XOR
g
r
R4
L4
T5

FIG.15

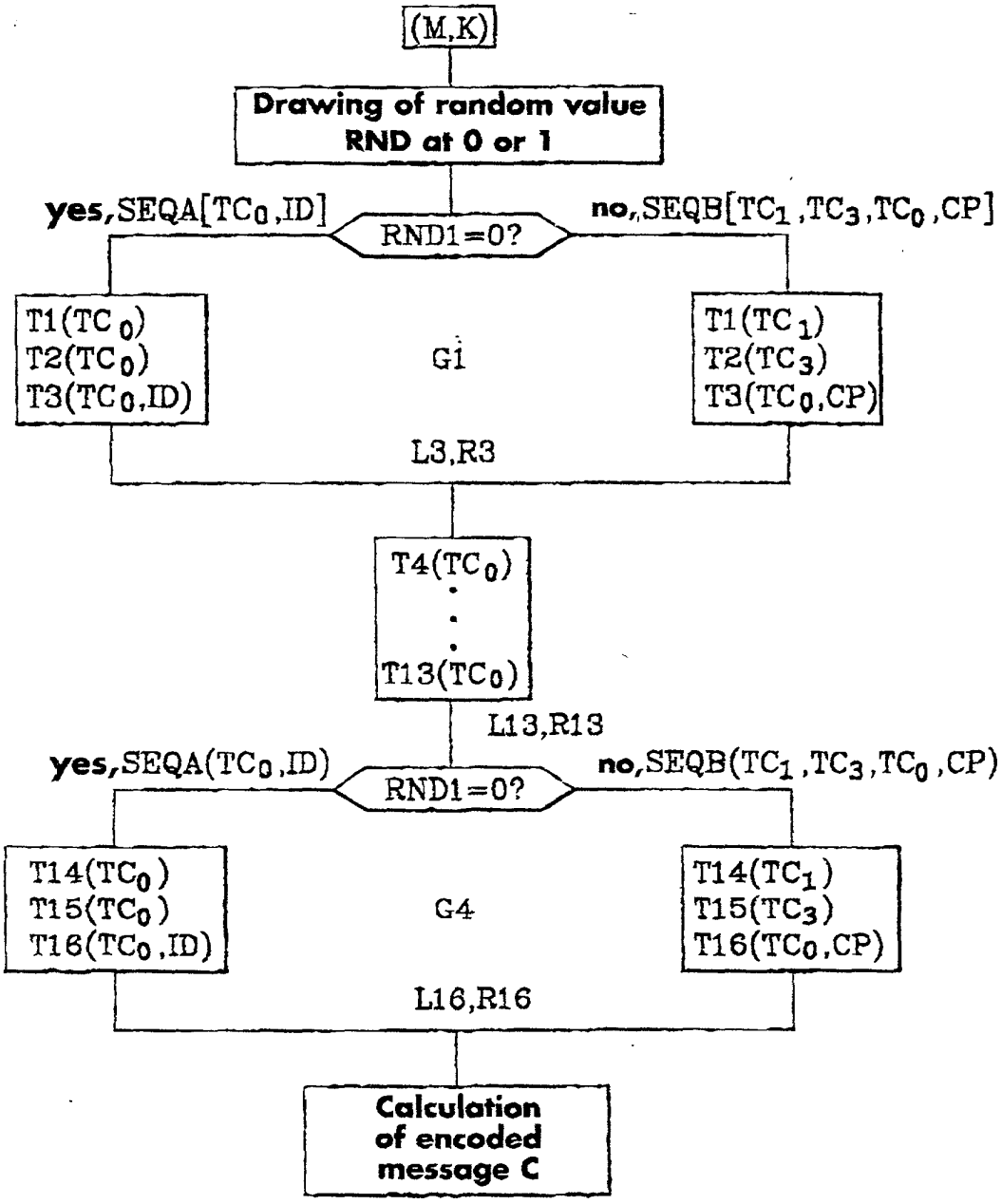

(M,K)
Drawing of random value RND at 0 or 1
yes,SEQA[TC0,ID]
RND1=0?
no,SEQB[TC1,TC3,TC0,CP]
T1(TC0)
T2(TC0)
T3(TC0,ID)
G1
T1(TC1)
T2(TC3)
T3(TC0,CP)
L3,R3
T4(TC0)
T13(TC0)
L13,R13
yes,SEQA(TC0,ID)
RND1=0?
no,SEQB(TC1,TC3,TC0,CP)
T14(TC0)
T15(TC0)
T16(TC0,ID)
G4
T14(TC1)
T15(TC3)
T16(TC0,CP)
L16,R16
Calculation of encoded message C

FIG.17

COMBINED DECLARATION FOR PATENT APPLICATION AND POWER OF ATTORNEY (Includes Reference to Provisional and International (PCT) Applications)	Attorney's Docket No. 032326-133

As a below named inventor, I hereby declare that:
My residence, post office address and citizenship are as stated below next to my name;
I BELIEVE I AM THE ORIGINAL, FIRST AND SOLE INVENTOR (IF ONLY ONE NAME IS LISTED BELOW) OR AN ORIGINAL, FIRST AND JOINT INVENTOR (IF PLURAL NAMES ARE LISTED BELOW) OF THE SUBJECT MATTER WHICH IS CLAIMED AND FOR WHICH A PATENT IS SOUGHT ON THE INVENTION ENTITLED:

Countermeasure Method in an Electronic Component Using a Secret Key Cryptographic Algorithm

The specification of which (check only one item below):

☐ is attached hereto.
☐ was filed as United States Patent Application Number 09/807,615
on April 16, 2001
and was amended on _____ (if applicable).

☐ was filed as International (PCT) Application Number _____
on _____
and was amended on _____ (if applicable).

I HAVE REVIEWED AND UNDERSTAND THE CONTENTS OF THE ABOVE-IDENTIFIED SPECIFICATION, INCLUDING THE CLAIMS, AS AMENDED BY ANY AMENDMENT REFERRED TO ABOVE.

I ACKNOWLEDGE THE DUTY TO DISCLOSE TO THE U.S. PATENT AND TRADEMARK OFFICE ALL INFORMATION KNOWN TO ME TO BE MATERIAL TO PATENTABILITY AS DEFINED IN TITLE 37, CODE OF FEDERAL REGULATIONS, Sec. 1.56 (as amended effective March 16, 1992);

I do not know and do not believe the said invention was ever known or used in the United States of America before my or our invention thereof, or patented or described in any printed publication in any country before my or our invention thereof or more than one year prior to said application; that said invention was not in public use or on sale in the United States of America more than one year prior to said application; that said invention has not been patented or made the subject of an inventor's certificate issued before the date of said application in any country foreign to the United States of America on any application filed by me or my legal representatives or assigns more than six months prior to said application;

I hereby claim foreign priority benefits under Title 35, United States Code, §§ 119 (a)-(e) of any foreign application(s) for patent or inventor's certificate or of any International (PCT) Application(s) designating at least one country other than the United States of America listed below and have also identified below any foreign application(s) for patent or inventor's certificate or any PCT International (PCT) Application(s) designating at least one country other than the United States of America filed by me on the same subject matter having a filing date before that of the application(s) of which priority is claimed:

PRIOR FOREIGN/PCT APPLICATION(S) AND ANY PRIORITY CLAIMS UNDER 35 U.S.C. §119:

COUNTRY (if PCT, indicate "PCT")	APPLICATION NUMBER	DATE OF FILING (day, month, year)	PRIORITY CLAIMED UNDER 35 U.S.C. §119
France	98/12990	16 October 1998	☐Yes ☐No

I hereby claim the benefit under Title 35, United States Code § 119(e) of any United States provisional application(s) listed below.

_____ _____
(APPLICATION NUMBER) (FILING DATE)

_____ _____
(APPLICATION NUMBER) (FILING DATE)

COMBINED DECLARATION FOR PATENT APPLICATION AND POWER OF ATTORNEY (CONT'D) (Includes Reference to Provisional and International (PCT) Applications)	Attorney's Docket No. 032326-133

I hereby claim the benefit under Title 35, United States Code, § 120 of any United States applications(s) or International (PCT) Application(s) designating the United States of America that is/are listed below and, insofar as the subject matter of each of the claims of this application is not disclosed in that/those prior application(s) in the manner provided by the first paragraph of Title 35, United States Code, § 112, I acknowledge the duty to disclose to the U.S. Patent and Trademark Office all information known to me to be material to the patentability as defined in Title 37, Code of Federal Regulations § 1.56, which became available between the filing date of the prior application(s) and the national or international filing date of this application:

PRIOR U.S. APPLICATIONS OR INTERNATIONAL (PCT) APPLICATIONS DESIGNATING THE U.S. FOR BENEFIT UNDER 35 U.S.C. § 120:

U.S. APPLICATIONS		STATUS (check one)		
U.S. APPLICATION NUMBER	U.S. FILING DATE	PATENTED	PENDING	ABANDONED
		☐	☐	☐

PCT APPLICATIONS DESIGNATING THE U.S.					
PCT APPLICATION NO.	PCT FILING DATE	U.S. APPLICATION NUMBERS ASSIGNED (if any)			
FR99/02199	15 September 1999				

I hereby appoint the following attorneys and agent(s) to prosecute said application and to transact all business in the U.S. Patent and Trademark Office connected therewith and to file, prosecute and to transact all business in connection with international applications directed to said invention:

William L. Mathis	17,337	Eric H. Weisblatt	30,505	Bruce T. Wieder	33,815
Robert S. Swecker	19,885	James W. Peterson	26,057	Todd R. Walters	34,040
Platon N. Mandros	22,124	Teresa Stanek Rea	30,427	Ronni S. Jillions	31,979
Benton S. Duffett, Jr.	22,030	Robert E. Krebs	25,885	Harold R. Brown III	36,341
Norman H. Stepno	22,716	William C. Rowland	30,888	Allen R. Baum	36,086
Ronald L. Grudziecki	24,970	T. Gene Dillahunty	25,423	Steven M. duBois	35,023
Frederick G. Michaud, Jr.	26,003	Patrick C. Keane	32,858	Brian P. O'Shaughnessy	32,747
Alan E. Kopecki	25,813	B. Jefferson Boggs, Jr.	32,344	Kenneth B. Leffler	36,075
Regis E. Slutter	26,999	William H. Benz	25,952	Fred W. Hathaway	32,236
Samuel C. Miller, III	27,360	Peter K. Skiff	31,917	Wendi L. Weinstein	34,456
Robert G. Mukai	28,531	Richard J. McGrath	29,195	Mary Ann Dillahunty	34,576
George A. Hovanec, Jr.	28,223	Matthew L. Schneider	32,814		
James A. LaBarre	28,632	Michael G. Savage	32,596		
E. Joseph Gess	28,510	Gerald F. Swiss	30,113		
R. Danny Huntington	27,903	Charles F. Wieland III	33,096		

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and: _____.

Address all correspondence to: James A. LaBarre
BURNS, DOANE, SWECKER & MATHIS, L.L.P.
P.O. Box 1404
Alexandria, Virginia 22313-1404

21839

Address all telephone calls to: James A. LaBarre at (703) 836-6620.

I hereby declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under Section 1001 of Title 18 of the United States Code and that such willful false statements may jeopardize the validity of the application or any patent issued thereon.

COMBINED DECLARATION FOR PATENT APPLICATION AND POWER OF ATTORNEY (CONT'D) (Includes Reference to Provisional and International (PCT) Applications)	Attorney's Docket No. 032326-133

FULL NAME OF SOLE OR FIRST INVENTOR Christophe CLAVIER	SIGNATURE	DATE 22/05/2001
RESIDENCE (CITY & STATE/COUNTRY) 5 rue de la Republique, F-13420 Gemenos, FRANCE		CITIZENSHIP France FRX
POST OFFICE ADDRESS (HOME ADDRESS) 5 rue de la Republique, F-13420 Gemenos, FRANCE		
FULL NAME OF SECOND JOINT INVENTOR, IF ANY Olivier BENOIT	SIGNATURE	DATE 11/05/2001
RESIDENCE (CITY & STATE/COUNTRY) La Treille d'Azur, Batiment D, avenue 19 Mars 1962, F-13400 Aubagne, FRANCE		CITIZENSHIP France FRX
POST OFFICE ADDRESS (HOME ADDRESS La Treille d'Azur, Batiment D, avenue 19 Mars 1962, F-13400 Aubagne, FRANCE		
FULL NAME OF THIRD JOINT INVENTOR, IF ANY	SIGNATURE	DATE
RESIDENCE (CITY & STATE/COUNTRY)		CITIZENSHIP
POST OFFICE ADDRESS (HOME ADDRESS		
FULL NAME OF FOURTH JOINT INVENTOR, IF ANY	SIGNATURE	DATE
RESIDENCE (CITY & STATE/COUNTRY)		CITIZENSHIP
POST OFFICE ADDRESS (HOME ADDRESS		
FULL NAME OF FIFTH JOINT INVENTOR, IF ANY	SIGNATURE	DATE
RESIDENCE (CITY & STATE/COUNTRY)		CITIZENSHIP
POST OFFICE ADDRESS (HOME ADDRESS		
FULL NAME OF SIXTH JOINT INVENTOR, IF ANY	SIGNATURE	DATE
RESIDENCE (CITY & STATE/COUNTRY)		CITIZENSHIP
POST OFFICE ADDRESS (HOME ADDRESS		
FULL NAME OF SEVENTH JOINT INVENTOR, IF ANY	SIGNATURE	DATE
RESIDENCE (CITY & STATE/COUNTRY)		CITIZENSHIP
POST OFFICE ADDRESS (HOME ADDRESS		
FULL NAME OF EIGHTH JOINT INVENTOR, IF ANY	SIGNATURE	DATE
RESIDENCE (CITY & STATE/COUNTRY)		CITIZENSHIP
POST OFFICE ADDRESS (HOME ADDRESS		
FULL NAME OF NINTH JOINT INVENTOR, IF ANY	SIGNATURE	DATE
RESIDENCE (CITY & STATE/COUNTRY)		CITIZENSHIP
POST OFFICE ADDRESS (HOME ADDRESS)		
FULL NAME OF TENTH JOINT INVENTOR, IF ANY	SIGNATURE	DATE
RESIDENCE (CITY & STATE/COUNTRY)		CITIZENSHIP
POST OFFICE ADDRESS (HOME ADDRESS		